

Exercise 12.1

Solve the following Linear Programming Problems graphically: 1. Maximise $Z = 3x + 4y$ subject to the constraints: $x + y \leq 4$, $x \geq 0$, $y \geq 0$.

Sol. Maximise $Z = 3x + 4y$... (i) subject to the constraints:

$x + y \leq 4$... (ii)

$x \geq 0$, $y \geq 0$... (iii) Step I. Constraint (iii) namely $x \geq 0$, $y \geq 0 \Rightarrow$ Feasible region

is in first quadrant. γ

Table of values for line $x + y =$

4

$\begin{matrix} 1 \\ + = \\ y \end{matrix} 4$

$B(0, 4)$

corresponding to constraint

(ii)

x

3

$x \ 0 \ 4$

2

y'

gives us $0 \leq 4$ which is true. Therefore region for constraint (ii) is on the origin side of the line.

The shaded region in the figure is the feasible region determined by the system of constraints (ii) and (iii). The feasible region OAB is bounded.

Step II. The coordinates of the corner points O, A and B are $(0, 0)$, $(4, 0)$ and $(0, 4)$ respectively.

Step III. Now we evaluate Z at each corner point.

Corner Point $Z = 3x + 4y$

$O(0, 0) \ 0$

$A(4, 0) \ 12$

$B(0, 4) \ 16 = M \leftarrow$ Maximum Hence, by Corner Point Method, the maximum value of Z is 16 attained at the corner point $B(0, 4)$. \Rightarrow Maximum $Z = 16$ at $(0, 4)$.

2. Minimise $Z = -3x + 4y$

subject to $x + 2y \leq 8$, $3x + 2y \leq 12$, $x \geq 0$, $y \geq 0$.

Sol. Minimise $Z = -3x + 4y$... (i) subject to: $x + 2y \leq 8$... (ii), $3x + 2y \leq 12$... (iii), $x \geq 0$, $y \geq 0$... (iv) Step I. Constraint (iv) namely $x \geq 0$, $y \geq 0 \Rightarrow$

Feasible region is in first quadrant.

Table of values for line $x + 2y = 8$ of constraint (ii)

$$\begin{array}{l} x \ 0 \ 8 \\ y \ 4 \ 0 \end{array}$$

Let us draw the line joining the constraint (ii) which gives $0 \leq 8$ which is true.

7 6 5

points (0, 4) and (8, 0). Now (ii) is \therefore Region let (0, 6) for us test for origin (0, 0) in 3

Y
1
on the origin side of the line. = 12 of constraint (iii) A(4, 0)(8, 0)
x
2
Table of values for line $3x + 2y = 8$
+ 2y = 8 y

points (0, 6) and (4, 0).
Let us draw the line joining the points (0, 6) and (4, 0).
X' O Y'

x 0 4
y 6 0
the line joining the points (0, 6) and (4, 0).
12 3 4 5
 $3 + 2 = 12$ x

Now let us test for origin (0, 0) in constraint (iii) which gives $0 \leq 12$ and which is true.

\therefore Region for constraint (iii) is also on the origin side of the line. The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (iv). The feasible region OABC is bounded.

Step II. The coordinates of the corner points O, A and C are (0, 0), (4, 0) and (0, 4) respectively.

Now let us find corner point B, intersection of lines $x + 2y = 8$ and $3x + 2y = 12$

Subtracting $2x = 4 \Rightarrow x = 2$.

Putting $x = 2$ in first equation $2 + 2y = 8$

$$\Rightarrow 2y = 6 \Rightarrow y = 3$$

\therefore Corner point B is (2, 3)

Step III. Now let us evaluate Z at each corner point.

$$\text{Corner Point } Z = -3x + 4y$$

$$O(0, 0) \quad 0$$

$$A(4, 0) - 12 = m \leftarrow \text{Minimum } B(2, 3) \quad 6$$

$$C(0, 4) \quad 16$$

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Hence, by Corner Point Method, the minimum value of Z is - 12 attained at the point A(4, 0).

\Rightarrow Minimum Z = - 12 at (4, 0).

3. Maximise $Z = 5x + 3y$

subject to $3x + 5y \leq 15$, $5x + 2y \leq 10$, $x \geq 0$, $y \geq 0$.

Sol. Maximise $Z = 5x + 3y$... (i) subject to:

$3x + 5y \leq 15$... (ii) $5x + 2y \leq 10$... (iii) $x \geq 0$, $y \geq 0$... (iv) Step

I. Constraint (iv) namely $x \geq 0$ and $y \geq 0$

\Rightarrow Feasible region is in first quadrant.

Table of values for line $3x + 5y = 15$ of constraint (ii) $x \geq 0$

Let us draw the line joining the points (0, 3) and (5, 0).
Let us test for origin (0, 0) in constraint (ii) which gives $0 \leq 15$ and which is true.

\therefore Region for constraint (ii)

$$C(0, 3) \quad B(5, 0)$$

contains the origin.

Table of values for line $5x + 2y = 10$ of constraint (iii).

2

$$3 + 5 = 15 \quad 0).$$

$$A(2, 0) \quad (5, 0) \quad X$$

$$12345^{5+2=1}_0x_67$$

Let us draw the line
joining the
points (0, 5) and (2,

Y'

Let us test for origin (0, 0) in constraint (iii) which gives $0 \leq 10$ and which is true.

\therefore Region for constraint (iii) also contains the origin. The shaded region in the figure is the feasible region determined by the system of constraints from (ii) and (iv). The feasible region OABC is bounded.

Step II. The coordinates of the corner points O, A and C are (0, 0), (2, 0) and (0, 3) respectively.

Now let us find corner point B; intersection of lines $3x + 5y = 15$ and $5x + 2y = 10$

$$\text{Ist eqn.} \times 2 - \text{IInd eqn.} \times 5 \text{ gives } -19x = -20 \Rightarrow x =$$

Putting $x =$

$$\text{in first eqn.} \Rightarrow +5y = 15$$

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$$\Rightarrow y =$$

. Therefore

$$\Rightarrow 5y = 15 - =$$

corner point B

$$\square \square \square$$

$$\square \square .$$

$$- \\ =$$

Step III. Now we evaluate Z at each corner point. Corner Point $Z = 5x + 3y$

$$O(0, 0) \quad 0$$

$$A(2, 0) \quad 10$$

+

$$B \quad =$$

$$\square \square \square$$

$$\square \square$$

$$C(0, 3) \quad 9$$

$$= M \leftarrow \text{Maximum}$$

Hence, by Corner Point Method, the maximum value of Z is attained at the corner point B

$$\begin{pmatrix} 3 \\ 0 \end{pmatrix}$$

$$\Rightarrow \text{Maximum } Z = 3x + 5y \text{ at } \begin{pmatrix} 3 \\ 0 \end{pmatrix}$$

4. Minimise Z =

such that $x + 3y \geq 3$, $x + y \geq 2$, $x, y \geq 0$.

Sol. Minimise $Z = 3x + 5y$... (i) such that: $x + 3y \geq 3$... (ii), $x + y \geq 2$... (iii), $x, y \geq 0$... (iv) Step I. The constraint (iv) $x, y \geq 0 \Rightarrow$ Feasible region is in first quadrant.

Table of values for line $x + 3y = 3$ of constraint (ii)

x	y
3	0
0	1

Let us draw the line joining the points (0, 1) and (3, 0). Now let us test for origin ($x = 0, y = 0$) in constraint (ii) $x + 3y \geq 3$, which gives us $0 \geq 3$ and which is not true.

\therefore Region for constraint (ii) does not contain the origin i.e., the region for constraint (ii) is not the origin side of the line. Table of values for line $x + y = 2$ of constraint (iii)

x 0 2

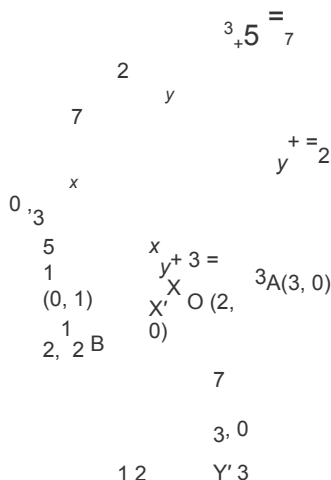
y 2 0

Let us draw the line joining the points (0, 2) and (2, 0). Now let us test for origin ($x = 0, y = 0$) in constraint (iii), $x + y \geq 2$, which gives us $0 \geq 2$ and which is not true. \therefore Region for constraint (iii) does not contain the origin i.e., is not the origin side of the line.

The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (iv). The feasible region is unbounded.

Step II. The coordinates of the corner points A and C are (3, 0) and (0, 2) respectively.

Y
C(0, 2)



Now let us find corner point B, the point of intersection of lines $x + 3y = 3$ and $x + y = 2$

Subtracting, $2y = 1 \Rightarrow y = \frac{1}{2}$

Putting $y = \frac{1}{2}$ in $x + y = 2$, we have $x = 2 - y = 2 - \frac{1}{2} = \frac{3}{2} \therefore$ Corner point

B is

$$\left(\frac{3}{2}, \frac{1}{2} \right)$$

Step III. Now, we evaluate Z at each corner point.

$$\text{Corner Point } Z = 3x + 5y$$

$$A(3, 0) \quad 9$$

B

$$\left(\frac{3}{2}, \frac{1}{2} \right) \quad \frac{13}{2}$$

$$+ = 7 = m \leftarrow \text{Smallest}$$

$$C(0, 2) \quad 10$$

From this table, we find that 7 is the smallest value of Z at the corner B

$\left(\frac{3}{2}, \frac{1}{2} \right)$. Since the feasible region is unbounded, 7 may or may not be the minimum value of Z.

Step IV. To decide this, we graph the inequality $Z < m$

i.e., $3x + 5y < 7$.

Table of values for line $3x + 5y = 7$
corresponding to constraint $3x + 5y < 7$ Class 12 Chapter

Let us draw the dotted
line joining the

x 0 y 0

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points

$\square \square \square$

$\square \square$ and

$\square \square \square$

$\square \square$. This line is to be shown dotted as

constraint involves $<$ and not \leq , so boundary of line is to be excluded. Let us test for origin ($x = 0, y = 0$) in constraint $3x + 5y < 7$, we have $0 < 7$ which is true. Therefore region for this constraint is on the origin side of the line $3x + 5y = 7$.

We observe that the half-plane determined by $Z < m$ has no point in common with the feasible region. Hence $m = 7$ is

the minimum value of Z attained at the point B

$\square \square \square$

$\square \square$.

\Rightarrow Minimum $Z = 7$ at

$\square \square \square$

$\square \square$.

5. Maximise $Z = 3x + 2y$

subject to $x + 2y \leq 10, 3x + y \leq 15, x, y \geq 0$.

Sol. Maximise $Z = 3x + 2y$... (i) subject to:

$x + 2y \leq 10$... (ii), $3x + y \leq 15$... (iii), $x, y \geq 0$... (iv) Step I. Constraint (iv) $x, y \geq 0 \Rightarrow$ Feasible region is in first quadrant.

Table of values for the line $x + 2y = 10$ corresponding to constraint (ii) x 0 10

y 5 0

Let us draw the line joining the points (0, 5) and (10, 0). Let us test for origin ($x = 0, y = 0$) in constraint (ii), we have $0 \leq 10$ which is true. \therefore Region for constraint (ii) is on the origin side of this line.

Table of values for line $3x + y = 15$ corresponding to constraint (iii)

$$x \leq 5$$

$$y \leq 15$$

Let us draw the line joining the points (0, 15) and (5, 0). Let us test for origin (x = 0, y = 0) in constraint (iii),

Region for we have $0 \leq 15$ which is true. \therefore

constraint (iii) is

and (5, 0). Let us test for origin (x = 0, y = 0) in constraint (iii),

$$5(0, 15)$$

The shaded region in the C(0, 5) B(4, 3)

line. on the origin side of this γ

region by the system of $x + 2y = 10$

figure is the feasible region determined

(ii) to (iv). γ'

OABC is bounded.

constraints from

The feasible region

$$y = 15 - x$$

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Step II. The coordinates of the corner points O, A and C are (0, 0), (5, 0) and (0, 5) respectively.

Now let us find corner point B, intersection of the lines $x + 2y = 10$ and $3x + y = 15$

First equation $- 2 \times$ second equation gives

$$-5x = 10 - 30 \Rightarrow -5x = -20 \Rightarrow x = 4$$

Putting $x = 4$ in $x + 2y = 10$, we have

$$4 + 2y = 10 \Rightarrow 2y = 6 \Rightarrow y = 3$$

\therefore Corner point B is B(4, 3).

Step III. Now we evaluate Z at each corner point.

Corner Point $Z = 3x + 2y$

$$O(0, 0) \quad 0$$

$$A(5, 0) \quad 15$$

$$B(4, 3) \quad 18 = M \leftarrow \text{Maximum}$$

C(0, 5) 10

Hence, by Corner Point Method, the maximum value of Z is 18 attained at the point B(4, 3).

⇒ Maximum Z = 18 at (4, 3).

6. Minimise $Z = x + 2y$

subject to $2x + y \geq 3$, $x + 2y \geq 6$, $x, y \geq 0$.

Show that the minimum of Z occurs at more than two points. Sol.

Minimise $Z = x + 2y$... (i) subject to:

$2x + y \geq 3$... (ii), $x + 2y \geq 6$... (iii), $x, y \geq 0$... (iv) Step I. Constraint

(iv) $x, y \geq 0 \Rightarrow$ Feasible region is in first quadrant.

Table of values for the line $2x + y = 3$ corresponding to constraint (ii).

x 0

y 3 0

Let us draw the line joining the points (0, 3) and (1.5, 0).

Now let us test for origin ($x = 0, y = 0$) in constraint (ii) $2x + y \geq 3$, we have $0 \geq 3$ which is not true.

∴ The region of constraint (ii) is on that side of the line which does not contain the origin i.e., the region other than the origin side of the line. Table of values for the line $x + 2y = 6$ corresponding to constraint (ii). x 0 6

y 3 0

Let us draw the line joining the points (0, 3) and (6, 0).

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Now let us test for origin ($x = 0, y = 0$) in constraint (iii) $x + 2y \geq 6$, we have $0 \geq 6$ which is not true.

∴ Region for constraint (iii) is the region other than the origin side of the line i.e., region not containing the origin. The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (iv). The feasible region is unbounded.

Step II. The coordinates of the corner points A and B are (6, 0) and (0, 3) respectively.

Y

$$B(0, 3) \quad 2x + 2y = 6 \quad y$$

$$x \quad x + 2y = 6$$

$$x + y + 2 < 6$$

$$x + y + 2 < 6$$

$$x + y + 2 < 6$$

$$x + y = 3$$

$$x + y = 3$$

$$x + y = 3$$

Step III. Now, we evaluate Z at each corner point.

Corner Point Z = x + 2y

A(6, 0) 6

= m ← Smallest

B(0, 3) 6

From this table, we find that 6 is the smallest value of Z at each of the two corner points. Since the feasible region is unbounded, 6 may or may not be the minimum value of Z. Step IV. To decide this, we graph the inequality $Z < m$ i.e., $x + 2y < 6$.

The line $x + 2y = 6$ for this constraint $Z < m$ ($\Rightarrow x + 2y < 6$) is the

same as the line AB for constraint (iii).

Let us test for origin ($x = 0, y = 0$) for this constraint, we have $0 < 6$ which is true.

Therefore region for this constraint is the (half-plane on) origin side of this line.

Points on the line segment AB are included in the feasible region and not in the half-plane determined by $x + 2y < 6$. We observe that the half-plane determined by $Z < m$ has no point in common with the feasible region. Hence $m = 6$ is the minimum

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value of Z attained at each of the points $A(6, 0)$ and $B(0, 3)$. \Rightarrow

Minimum $Z = 6$ at $(6, 0)$ and $(0, 3)$.

Remark. In fact, $Z = 6$ at all points on the line segment AB for example

$(2, 2)$, etc.

7. Minimise and Maximise $Z = 5x + 10y$ subject to $x + 2y \leq 120$, $x + y \geq 60$, $x - 2y \geq 0$, $x, y \geq 0$.

Sol. Minimise and Maximise $Z = 5x + 10y$... (i) subject to: $x + 2y \leq 120$... (ii) $x + y \geq 60$... (iii), $x - 2y \geq 0$... (iv), $x, y \geq 0$... (v) Step I. Constraint (v) $x, y \geq 0 \Rightarrow$

Feasible region is in first quadrant. Table of values for line $x + 2y = 120$ of constraint (ii) $x \geq 0$

$y \geq 0$

Let us draw the line joining the points $(0, 60)$ and $(120, 0)$. Let us test for origin ($x = 0, y = 0$) in constraint (iii) $x + 2y \leq 120$ we have $0 \leq 120$ which is true.

\therefore Region for constraint (ii) is on the origin side of the line $x + 2y = 120$.

Table of values for line $x + y = 60$ of constraint (iii)

$x \geq 0$

$y \geq 0$

Let us draw the line joining the points $(0, 60)$ and $(60, 0)$. Let us test for origin ($x = 0, y = 0$) in constraint (iii) $x + y \geq 60$, we have $0 \geq 60$ which is not true.

\therefore Region for constraint (iii) is the half-plane on the non-origin side of the line $x + y = 60$ (i.e., on the side of the line opposite to the origin side).

Table of values for line $x - 2y = 0$ of constraint (iv)

$x \geq 0$

$y \geq 0$

The line $x - 2y = 0$ is passing through the origin, so we still have another point $(60, 30)$

on the line). Let us test for (60, 0)

Let us draw the line $x + y = 60$

joining the points (0, 0) and (60, 30).

(a point other than origin) in constraint (iv), we have $60 \geq 0$ which is true.

\therefore Region for (iv) is the half-plane on that side of the line which containing the point (60, 0).

x 20 40 80 100

y 0 20 30 40 60

$x - 2y = 0$

$x + y = 60$

D(40, 20) C(60, 30)

A(60, 0) B(120, 0)

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The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (v). The feasible region ABCD is bounded.

Step II. The coordinates of the corner points A and B are (60, 0) and (120, 0) respectively.

Corner point C is the intersection of the line $x - 2y = 0$ i.e., $x = 2y$ and $x + 2y = 120$. Putting $x = 2y$ in $x + 2y = 120$, we have $2y + 2y = 120 \Rightarrow 4y = 120 \Rightarrow y = 30$ and therefore $x = 2y = 60$.

\therefore Corner point C (60, 30).

Similarly for corner point D, putting $x = 2y$ in $x + y = 60$, we have $2y + y = 60 \Rightarrow 3y = 60 \Rightarrow y = 20$ and therefore $x = 2y = 40$. Therefore corner point D is (40, 20).

Step III. Now, we evaluate Z at each corner point.

$$\text{Corner Point } Z = 5x + 10y$$

$$A(60, 0) \quad 300 = m \leftarrow \text{Minimum}$$

$$B(120, 0) \quad 600$$

$$C(60, 30) \quad 300 + 300 = 600 = M \leftarrow \text{Maximum} \quad D(40, 20) \quad 400$$

Hence, by Corner Point Method,

Minimum Z = 300 at (60, 0)

Maximum Z = 600 at B(120, 0) and C(60, 30) and hence maximum at all the points on the line segment BC joining the points (120, 0) and (60, 30).

8. Minimise and Maximise $Z = x + 2y$

subject to $x + 2y \geq 100$, $2x - y \leq 0$, $2x + y \leq 200$; $x, y \geq 0$. Sol.

Minimise and Maximise $Z = x + 2y$... (i) subject to:

$$x + 2y \geq 100 \dots(ii)$$

$$2x - y \leq 0 \dots(iii)$$

$$2x + y \leq 200 \dots(iv)$$

$$x, y \geq 0 \dots(v)$$

Step I. The constraint (v) $x, y \geq 0 \Rightarrow$ Feasible region is in first quadrant. Table of values for the line $x + 2y = 100$ for constraint (ii). x

$$y \ 50 \ 0$$

Let us draw the line joining the points (0, 50) and (100, 0). Let us test for origin ($x = 0, y = 0$) in constraint (ii) $x + 2y \geq 100$, we have $0 \geq 100$ which is not true.

\therefore Region for constraint (i) is that half-plane which does not contain the origin.

Table of values for the line $2x - y = 0$ i.e., $2x = y$ of constraint (iii).

$$x \ 0 \ 20$$

$$y \ 0 \ 40$$

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Let
us

draw the line joining the points (0, 0) and (20, 40). Because this line passes through the origin, so we shall have the test for some point say (100, 0) other than the origin. Putting $x = 100$ and $y = 0$ in constraint (iii) $2x - y \leq 0$, we have $200 \leq 0$ which is not true.

\therefore Region for constraint (iii) is the half plane on the side of the line which does not contain the point (100, 0).

Table of values for the line $2x + y = 200$ of constraint (iv).

$$x \ 0 \ 100$$

$$y \ 200 \ 0$$

Let us draw the line joining the points (0, 200) and (100, 0). Let us test for origin ($x = 0, y = 0$) in constraint (iv) $2x + y \leq 200$, we have $0 \leq 200$ which is true. Therefore region for constraint (iv) is the half-plane on origin side of the line.

The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (v). The feasible region ABCD is bounded.

Step II. The C(0, 200) is the $-y = 0$ i.e., y
coordinates and D(0, intersection = $2x$.
of the two 50). of boundary
corner Corner lines $x + 2y$
points are point A = 100 and $2x^{200}$

$$C(0, 200)$$

$$= 0 \quad y$$

$$2 - x$$

150

putting $y = 2x$, x

$$+ 4x = 100 \Rightarrow 5x \quad D(0, 50)$$

$$= 100 \Rightarrow x = 20. \quad B(50, 100)$$

$$\therefore y = 2x = 2 \times 20 \quad A(20, 40)$$

$$= 40. \quad (100, 0)$$

Solving them,

Therefore point B is boundary y , $2 + 150x$
corner the lines $2x + x \leq 50$ x

$$\text{point } A(20, 40). = 200 \quad y + 2x = 200 \Rightarrow 4x \quad y + 2 = 100$$

$$\text{Corner } y' \text{ of the } + 2x = 200 \Rightarrow 4x \quad y = 200$$

$$\text{intersection } x = 200$$

$$\text{and } 2x - y = 0 \text{ i.e., } = 200$$

$$y = 2x. \text{ Solving}$$

$$\text{them, putting } y =$$

$\Rightarrow x = 50$ and therefore $y = 2x = 100$. Therefore corner point B is (50, 100).

Step III. Now, we evaluate Z at each corner point.

$$\text{Corner Point } Z = x + 2y$$

$$A(20, 40) \quad 100 = m \leftarrow \text{Minimum} \quad B(50, 100) \quad 250$$

$$C(0, 200) \quad 400 = M \leftarrow \text{Maximum} \quad D(0, 50) \quad 100 = m \leftarrow \text{Minimum}$$

By Corner Point Method,

Minimum $Z = 100$ at all the points on the line segment joining the points (20, 40) and (0, 50).

(See Step III, Example 7, Page 770.

Maximum $Z = 400$ at (0, 200).

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9. Maximise $Z = -x + 2y$, subject to the constraints: $x \geq 3$, $x + y \geq 5$, $x + 2y \geq 6$, $y \geq 0$.

Sol. Maximise $Z = -x + 2y$... (i) subject to the constraints: $x \geq 3$... (ii), $x + y \geq 5$... (iii), $x + 2y \geq 6$... (iv), $y \geq 0$... (v) Step I. Constraint (v), $y \geq 0 \Rightarrow$

Positive side of y-axis \Rightarrow Feasible region is in first and second quadrants. Region for constraint (ii) $x \geq 3$. We know that graph of the line $x = 3$ is a vertical line parallel to y-axis at a distance 3 from origin along OX.

\therefore Region for $x \geq 3$ is the half-plane on right side of the line $x = 3$. Table

of values for line $x + y = 5$ of constraint (iii) $x \geq 0$
 $y \geq 0$

Let us draw the line joining the points (0, 5) and (5, 0). Let us test for origin (0, 0) in constraint (ii).

Putting $x = 0$ and $y = 0$ in $x + y \geq 5$, we have $0 \geq 5$ which is not true. \therefore

Region for constraint (iii) is the half plane on the non-origin side of the line $x + y = 5$.

Table of values for the line $x + 2y = 6$ of constraint (iii) $x \geq 0$
 $y \geq 0$

Let us test for origin (0, 0) in constraint (iv) $x + 2y \geq 6$, we have $0 \geq 6$ which is not true.

\therefore Region for constraint

(iv) is again the half plane on the non-origin side of the line $x + 2y = 6$. The shaded region in the

region determined by the

the line $x + 2y =$

(0, 5) $x = 3$

$- + 2 > 1 \times y$
figure is the feasible

$+ = 5$
 y
(0, 3)

$- x + 2 = 1$

from (ii) to feasible (v).

The region $(-1, 0)$

system of constraint is

$0, 1, 2$

(6, 0).

Step II. The coordinates of the corner point A are

y'

Corner point B is the

$+ 2 = 6$ $x' O$

C(3, 2) B(4, 1)

unbounded.

1 2 3 4 5 X
7 X

intersection of the boundary lines

(5, 0)
A(6, 0)

$x + y = 5$ and $x + 2y = 6$

Let us solve them for x and y .

Subtracting the two equations $2y - y = 6 - 5$ or $y = 1$. Putting $y = 1$ in $x + y = 5$, we have $x + 1 = 5$ or $x = 4$. Therefore, vertex B is (4, 1). Corner point C is the intersection of the boundary lines $x + y = 5$ and $x = 3$.

2. Therefore corner point C is (3, 2).

Step III. Now, we evaluate Z at each corner point.

$$\text{Corner Point } Z = -x + 2y$$

$$A(6, 0) - 6$$

$$B(4, 1) - 2$$

$C(3, 2) 1 = M \leftarrow$ Maximum From this table, we find that 1 is the maximum value of Z at (3, 2). Step IV. Since the feasible region is unbounded, 1 may or may not be the maximum value of Z. To decide this, we graph the inequality $Z > M$ i.e., $-x + 2y > 1$.

Table of values for the line $-x + 2y = 1$ corresponding to constraint $Z > M$ i.e., $-x + 2y > 1$.

$$x \ 0 \ -1$$

y

$$\square \square \square$$

$$\square \square \text{ and } (-1, 0).$$

Let us draw the dotted line joining the points

The line is to be shown dotted because boundary of the line is to be excluded as equality sign is missing in the constraint $Z > M$. We observe that the half-plane determined by $Z > M$ has points in common with the feasible region. Therefore, $Z = -x + 2y$ has no maximum value subject to the given constraints.

10. Maximise $Z = x + y$,

subject to $x - y \leq -1$, $-x + y \leq 0$, $x, y \geq 0$.

Sol. Maximise $Z = x + y$... (i) subject to:

$x - y \leq -1$... (ii), $-x + y \leq 0$... (iii), $x, y \geq 0$... (iv) Step I.

Constraint (iv) $x, y \geq 0$.

\Rightarrow Feasible region is in first quadrant.

Table of values for the line $x - y = -1$ of constraint (ii) $x \ 0 \ -1$

$$y \ 1 \ 0$$

Let us draw the straight line joining the points (0, 1) and (-1, 0). Let us test for origin (0, 0) in constraint (ii) $x - y \leq -1$, we have $0 \leq -1$ which is not true.

Therefore region for constraint (ii) is the region on that side of the line which is away from the origin (as shown shaded in the figure) Table of values for the line $-x + y = 0$ i.e., $y = x$ of constraint (iii) $x \ 0 \ 2$

$$y \ 0 \ 2$$

Let us draw the line joining the points (0, 0) and (2, 2). Let us test for the point (2, 0) (say) [and not origin as line passes through (0, 0)] in

constraint (iii) $-x + y \leq 0$, we have $-2 \leq 0$ which is true.

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point

(2, 0) side of the line

\therefore Region for
constraint

in the

Y (iii) is towards (shown shaded

$y_- \leq$

-1

(2, 2)

figure). 2

From observe

$y_+ = 0 \quad y_- = x$

the that

common

$- \quad x$

x_-

in the 1

< 0

(2, 0)

figure, we no point x_+
there is two (0, 1) y_+

shaded regions. Thus, the

$x'x$

solution i.e., no problem has no 12
O

feasible region and

hence no feasible

maximum value of Z.

(-1, 0)

Y'

Exercise 12.2

1. Reshma wishes to mix two types of food P and Q in such a way that the vitamin contents of the mixture contain at least 8 units of vitamin A and 11 units of vitamin B. Food P costs ₹ 60/kg and Food Q costs ₹ 80/kg. Food P contains 3 units/kg of vitamin A and 5 units/kg of vitamin B while food Q contains 4 units/kg of vitamin A and 2 units/kg of vitamin B. Determine the minimum cost of the mixture.

Sol. Step I. Mathematical formulation of L.P.P.

Suppose Reshma mixes x kg of food P and y kg of food Q. The given data is condensed in the following table:

Type	Quantity	Cost	Vitamin A	Vitamin B	Food (kg)	(`/kg)
	(units/kg)	(units/kg)	(units/kg)	(units/kg)	P	Q
					60	80
					3	4
					5	2

Cost of mixture (in `) = $60x + 80y$

Let $Z = 60x + 80y$

We have the following mathematical model for the given problem:

Minimise $Z = 60x + 80y$... (i) subject to the constraints:

$3x + 4y \geq 8$ (Vitamin A constraint) ... (ii) [Given: Vitamin A content of foods X and Y is at least (i.e., \geq) 8 units]

$5x + 2y \geq 11$ (Vitamin B constraint) ... (iii) [Given: Vitamin B content of foods X and Y is at least (i.e., \geq) 11 units]

$x, y \geq 0$ [Quantities of food can't be negative] ... (iv) Step II. The constraint (iv), $x, y \geq 0$.

\Rightarrow Feasible region is in first quadrant.

Table of values for the line $3x + 4y = 8$ of constraint (ii)

y	x
2	0
0	8/3

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Let us draw the line joining the points (0, 2) and (8/3, 0).

Let us test for origin ($x = 0, y = 0$) in constraint (ii) $3x + 4y \geq 8$, we have $0 \geq 8$ which is not true.

\therefore The region for constraint (ii) is the half plane on non-origin side of the line $3x + 4y = 8$ i.e., the region does not contain the origin. Now

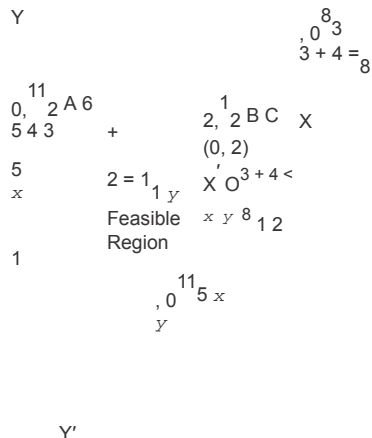
table of values for the line $5x + 2y = 11$ of constraint (iii).

y	x
11/2	0
0	11/5

Let us draw the line joining the points (0, 11/2) and (11/5, 0).

Let us test for origin ($x = 0, y = 0$) in constraint (iii) $5x + 2y \geq 11$, we have $0 \geq 11$ which is not true.

\therefore Region for constraint (iii) is on the non-origin side of the line i.e., does not contain the origin.



The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (iv). The feasible region is unbounded.

Step III. The coordinates of the corner points A and C are \square \square \square \square and C \square \square \square \square respectively.

Corner point B; is the point of intersection of the lines

$$3x + 4y = 8 \text{ and } 5x + 2y = 11$$

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$$3x + 4y = 8 \text{ and } 5x + 2y = 11$$

Solve for x and y: First equation $-2 \times$ second equation gives $3x + 4y - 10x - 4y = 8 - 22$

$$\Rightarrow -7x = -14 \Rightarrow x = 2$$

Putting $x = 2$ in $3x + 4y = 8$, we have, $6 + 4y = 8 \Rightarrow 4y = 2 \Rightarrow y = \frac{1}{2}$

Therefore vertex B

Corner Point Z = 60x + 80y

B

$\square \square \square \quad 160$
 $\square \square = m \leftarrow \text{Minimum}$

C

□ □ □ □ □ 160

the two corner points B

□ □ □
□ □ and C
□ □ □
□ □

Step V. Since the feasible region is unbounded, 160 may or may not be the minimum value of Z. To decide this, we graph the inequality $Z < m$ i.e., $60x + 80y < 160$ or $3x + 4y < 8$.

Table of values for the line $3x + 4y = 8$ for this constraint $Z < m$.

$$\begin{matrix} x & 0 \\ y & 2 & 0 \end{matrix}$$

The line joining these two points $(0, 2)$ and $(1, 1)$ has already

been drawn for the line of constraint (ii).

Let us test for origin ($x = 0, y = 0$) in constraint $Z < m$ i.e., $3x + 4y < 8$, we have $0 < 8$ which is true.

∴ Region for constraint $Z < m$ in the origin side of the line $3x + 4y = 8$.

Of course points on the line segment BC are included in the feasible

region (i) of constraint (ii) and not in the half-plane determined by $Z < m$ i.e., $3x + 4y < 8$. We observe that the half-plane determined by $Z < m$ has no point in common with the feasible region. Hence $m = 160$ is the minimum value of Z .

attained at each of the points B

$\frac{1}{2}$ and C
 $\frac{1}{2}$. Therefore,

minimum cost = ₹ 160 at all points lying on the segment joining

and

$\frac{1}{2}$. 17
 $\frac{1}{2}$.

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2. One kind of cake requires 200 g of flour and 25 g of fat, and another kind of cake requires 100 g of flour and 50 g of fat. Find the maximum number of cakes which can be made from 5 kg of flour and 1 kg of fat assuming that there is no shortage of the other ingredients used in making the cakes.

Sol. Step I. Mathematical Formulation of L.P.P.

Let x be the number of cakes of first kind and y , the number of cakes of other kind. The given data is condensed in the following table:

Kind of cake	Number of cakes	Flour (gm/cake)	Fat (gm/cake)
I	x	200	25
II	y	100	50

I x 200 25

II y 100 50

Total number of cakes = $x + y$ Let $Z = x + y$ We have the following mathematical model for the given problem: Maximise $Z = x + y$... (i) subject to the constraints:

$200x + 100y \leq 5000$

(Given: (Maximum) amount of flour available for both types of cakes is 5 kg = 5000 gm)

Dividing by 100,

or $2x + y \leq 50$ (Flour constraint) ... (ii) $25x + 50y \leq 1000$

(Fat constraint)

(Given: (Maximum) amount of fat available for both types of cakes is 1 kg = 1000 gm)

Dividing by 25,

or $x + 2y \leq 40$ (Fat constraint) ... (iii) $x, y \geq 0$... (iv)

(∵ Number of cakes can't be negative)

Step II. The constraint (iv) $x, y \geq 0$.

⇒ Feasible region is in first quadrant.

Table of values for the line $2x + y = 50$ of constraint (ii)

x	y
0	50
25	0

y 50 0

Let us draw the line joining the points (0, 50) and (25, 0). Let us test for origin (0, 0) ($x = 0$ and $y = 0$) in constraint (ii) $2x + y \leq 50$, we have $0 \leq 50$ which is true.

∴ Region for constraint (ii) is towards the origin side of the line.
 Table of values for the line $x + 2y = 40$ of constraint (iii) $x \geq 0, y \geq 0$
 Let us draw the line joining the points (0, 20) and (40, 0). Let us test for origin ($x = 0, y = 0$) in constraint (iii) $x + 2y \leq 40$, we have $0 \leq 40$ which is true.
 ∴ Region for constraint (iii) is also towards the origin side of the line.
 The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (iv). The feasible region is bounded.

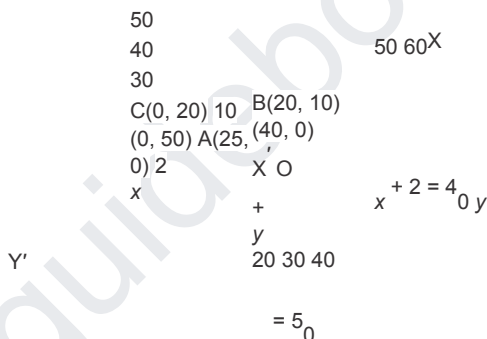
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Step III. The coordinates of the corner points O, A and C are (0, 0), (25, 0) and (0, 20) respectively.

Corner point B: It is the point of intersection of the boundary lines $2x + y = 50$ and $x + 2y = 40$

Y
10



Let us solve them for x and y.

First equation $- 2 \times$
 second equation
 gives

$$2x + y - 2x - 4y = 50 - 80 \Rightarrow -3y = -30 \Rightarrow y = 10. \text{ Putting } y = 10 \text{ in}$$

$$2x + y = 50$$

$$\Rightarrow 2x + 10 = 50 \Rightarrow 2x = 40 \Rightarrow x = 20$$

Therefore corner point B is (20, 10).

Step IV. Now we evaluate Z at each corner point.

Corner Point $Z = x + y$

$$O(0, 0) \quad 0$$

$$A(25, 0) \quad 25$$

$$B(20, 10) \quad 30 = M \leftarrow \text{Maximum}$$

$$C(0, 20) \quad 20$$

By Corner Point Method, the maximum value of Z is 30 attained at the

point B(20, 10).

Hence, maximum number of cakes = 30, 20 of first kind and 10 of second kind.

3. A factory makes tennis rackets and cricket bats. A tennis racket takes 1.5 hours of machine time and 3 hours of craftsman's time in its making while a cricket bat takes 3 hours of machine time and 1 hour of craftsman's time. In a day, the factory has the availability of not more than 42 hours of machine time and 24 hours of craftsman's time.

(i) What number of rackets and bats must be made if the factory is to work at full capacity?

(ii) If the profit on a racket and on a bat is ₹ 20 and ₹ 10 respectively, find the maximum profit of the factory when it works at full capacity.

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Sol. Step I. Mathematical Formulation of L.P.P.

Suppose x is the number of tennis rackets and y is the number of cricket bats to be made in a day. The given data is condensed in the following table:

Item	Number	Machine Time (hours/item)	Craftsman's Time (hours/item)	Profit (₹/item)
------	--------	---------------------------	-------------------------------	-----------------

Tennis Racket x 1.5 3 20 Cricket Bat y 3 1 10
 Total number of items = $x + y$ and total profit = $20x + 10y$
 Let $Z = x + y$ and $P = 20x + 10y$
 We have the following mathematical model for the given problem: Maximise $Z = x + y$ and $P = 20x + 10y$... (i) subject to the constraints:

$$1.5x + 3y \leq 42 \text{ or } x + 3y \leq 42$$

[Given: Number of machine hours available is not more than 42 hours i.e., ≤ 42]

Dividing by 3 and multiplying by 2,

$$x + 2y \leq 28 \text{ (Machine time constraint) ... (ii) } 3x + y \leq 24$$

(Craftsman's time constraint) ... (iii) [Given: Number of craftsman's hours is not more than 24 hours i.e., ≤ 24]

$$x, y \geq 0$$

(∵ Number of tennis rackets and cricket bats can't be negative) ... (iv)

Step II. The constraint (iv) $x, y \geq 0 \Rightarrow$ Feasible region is in first quadrant.

Table of values of equation $x + 2y = 28$ of constraint (ii)

x	0	28
y	14	0

Let us draw the straight line joining the points (0, 14) and (28, 0). Let us test for origin ($x = 0, y = 0$) in constraint (ii)

i.e., $x + 2y \leq 28$; we have $0 \leq 28$ which is true.

∴ Region for constraint (ii) is the region towards the origin side of the line $x + 2y = 28$.

Table of values of equation $3x + y = 24$ of constraint (iii) $x \leq 8$
 $y \leq 24$

Let us draw the line joining the points $(0, 24)$ and $(8, 0)$. Let us test for origin $(x = 0, y = 0)$ in constraint (iii) $3x + y \leq 24$, we have $0 \leq 24$ which is true.

\therefore Region for constraint (iii) is the region towards the origin side of the line.

The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (iv). The feasible region is bounded.

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Step III. The coordinates of the corner points O, A and C are $(0, 0)$, $(8, 0)$ and $(0, 14)$ respectively.

Corner point B: It is the point of intersection of the boundary lines $x + 2y = 28$ and $3x + y = 24$.

First eqn. $\times 2$ \times second eqn. gives

$$x + 2y - 2(3x + y) = 28 - 2 \times 24$$

$$\Rightarrow x + 2y - 6x - 2y = 28 - 48 \Rightarrow -5x = -20$$

$$\Rightarrow x = 4.$$

Putting $x = 4$ in $x + 2y = 28$, $4 + 2y = 28$

$$\Rightarrow 2y = 24 \Rightarrow y = 12$$

\therefore Corner point B is $(4, 12)$.

Y

$(0, 24)$

20

16

C(0, 14) A(8, 0) (28, 0)

12 8

B(4, 12) $\times 4$ 12 16 20 24

$$+ 2 = 28$$

4

x

+

X' O Y'

y

$$= 24$$

Step IV. (i) Now, we evaluate Z at each corner point.

Corner Point $Z = x + y$

$O(0, 0)$ 0

$A(8, 0)$ 8

$B(4, 12)$ 16 = $M \leftarrow$ Maximum $C(0, 14)$ 14

By Corner Point Method, maximum $Z = 16$ at $(4, 12)$. (ii)

Now, we evaluate P at each corner point.

Corner Point $P = 20x + 10y$

$O(0, 0)$ 0

$A(8, 0)$ 160

$B(4, 12)$ 200 = $M \leftarrow$ Maximum $C(0, 14)$ 140

By Corner Point Method, maximum $P = 200$ at $(4, 12)$. Hence, the factory should make 4 tennis rackets and 12 cricket bats to make use of full capacity and then the profit is also maximum equal to ` 200.

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4. A manufacturer produces nuts and bolts. It takes 1 hour of work on machine A and 3 hours on machine B to produce a package of nuts. It takes 3 hours on machine A and 1 hour on machine B to produce a package of bolts. He earns a profit of ` 17.50 per package on nuts and ` 7.00 per package on bolts. How many packages of each should be produced each day so as to maximise his profit, if he operates his machines for at the most 12 hours a day?

Sol. Sol. Step I. Mathematical Formulation of L.P.P. Suppose the manufacturer produces x packages of nuts and y packages of bolts each day. The given data is condensed in the following table:

Item	Number of hours per package on Machine A	Number of hours per package on Machine B	Profit (/package)
Nuts x	1	3	17.50
Bolts y	3	1	7.00

Total profit (in `) = $17.5x + 7y$

Let $Z = 17.5x + 7y$

We have the following mathematical model for the given problem.

Maximise $Z = 17.5x + 7y$... (i) subject to the constraints:

$x + 3y \leq 12$ (Machine A constraint) ... (ii) (Given: He operates his machine A for at most 12 hours i.e., ≤ 12 hours)

$3x + y \leq 12$ (Machine B constraint) ... (iii) (Given: He operates his machine B also for at the most 12 hours i.e., ≤ 12 hours)

$x, y \geq 0$... (iv) \therefore Number of packages of nuts and bolts can't be negative) Constraint (iv) $x, y \geq 0$

\Rightarrow Feasible region is in first quadrant.

Step-II. Table of values for the line $x + 3y = 12$ of constraint (ii)

x	y
0	4
12	0

Let us draw the straight line joining the points (0, 4) and (12, 0). Let us test for origin ($x = 0, y = 0$) in constraint (ii).

$x + 3y \leq 12$, we have $0 \leq 12$ which is true.

\therefore Region for constraint (ii) is the region on the origin side of the line $x + 3y = 12$.

Table of values for the line $3x + y = 12$ of constraint (iii)

x	y
0	12
4	0

Let us draw the straight line joining the points (0, 12) and (4, 0). Let us test for origin ($x = 0, y = 0$) in constraint (iii) $3x + y \leq 12$, we have $0 \leq 12$ which is true.

\therefore Region for constraint (iii) is also on the origin side of the line $3x + y = 12$.

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Y

(0, 12)

3+

x
8

= 12

$\frac{x}{y} + 3 = 12$

C(0, 4)

B(3, 3)

X O A(4, 0) Y

8 (12, 0)

The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (iv). The feasible region is bounded.

Step III. The coordinates of the corner points O, A and C are (0, 0), (4, 0) and (0, 4) respectively.

Corner point B: It is the point of intersection of the boundary lines $x + 3y = 12$ and $3x + y = 12$

Solving them for x, y:

1st Eqn. $- 3 \times$ second Eqn. gives

$$x + 3y - 3(3x + y) = 12 - 36$$

$$\Rightarrow x + 3y - 9x - 3y = -24 \Rightarrow -8x = -24$$

$$\Rightarrow x = 3$$

evaluate Z at each corner point.

$$\text{Corner Point } Z = 17.5x + 7y$$

$$O(0, 0) \quad 0$$

$$A(4, 0) \quad 70$$

$$\text{Putting } x = 3 \text{ in } x + 3y = 12, \quad 3 + 3y = 12$$

$$\Rightarrow 3y = 9 \Rightarrow y = 3$$

\therefore Corner point B is (3, 3).

Step IV. Now, we

$$B(3, 3) \quad 73.5 = M \leftarrow \text{Maximum } C(0, 4) \quad 28$$

By Corner Point Method, maximum Z = 73.5 at (3, 3). Hence, the profit is maximum equal to ` 73.50 when 3 packages of nuts and 3 packages of bolts are manufactured.

5. A factory manufactures two types of screws, A and B. Each type of screw requires the use of two machines, an automatic and a hand operated. It takes 4 minutes on the automatic and 6 minutes on hand operated machines to

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manufacture a package of screws A, while it takes 6 minutes on automatic and 3 minutes on the hand operated machines to manufacture a package of screws B. Each machine is available for at the most 4 hours on any day. The manufacturer can sell a package of screws A at a profit of ` 7 and screws B at a profit of ` 10.

Assuming that he can sell all the screws he manufactures, how many packages of each type should the factory owner produce in a day in order to maximise his profit? Determine the maximum profit.

Sol. Step I. Mathematical Formulation of L.P.P.

Suppose the factory owner produces x packages of screw A and y packages of screw B in a day. The given data is condensed in the following table:

Type of screw	Number of packages	Time in minutes per item	Profit (₹/item)
A	x	4	7
B	y	6	10
Total		Profit = 7x + 10y	

$$\text{Let } Z = 7x + 10y$$

We have the following mathematical model for the given problem.

Maximise $Z = 7x + 10y$... (i) subject to the constraints:

$$4x + 6y \leq 240$$

[Each machine i.e., automatic machine is also available for at most i.e., ≤ 4 hours i.e., $4 \times 60 = 240$ minutes]

or $2x + 3y \leq 120$ (Automatic machine constraint) ... (ii) $6x + 3y \leq 240$
 (Same argument as given above for constraint (ii))

or $2x + y \leq 80$... (iii) (Hand operated machine constraint)

$x, y \geq 0$... (iv) \therefore Number of screws A and B can't be negative)

Step II. Table of values for the line $2x + 3y = 120$ of constraint (ii)

x 0 60

y 40 0

Let us draw the straight line joining the points (0, 40) and (60, 0). Let us test for origin (put $x = 0, y = 0$) in constraint (ii) $2x + 3y \leq 120$, we have $0 \leq 120$ which is true.

\therefore Region for constraint (ii) is on the origin side of the line $2x + 3y = 120$.

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Table of values for the line $2x + y = 80$ of constraint (iii)

x 0 40

y 80 0

Let us draw the straight line joining the points (0, 80) and (40, 0). Let us test for origin (put $x = 0, y = 0$) in constraint (iii) $2x + y \leq 80$, we have $0 \leq 80$ which is true.

\therefore Region for constraint (iii) is also towards the origin side of the line $2x + y = 80$.

Y

80 (0, 80)

70 60 50 40 30 20

10

2

x

+

y

=

80

C(0, 40)

B(30, 20)

(60, 0)

X' O Y'

10 20 30

50 A(40,

0)

60 70

2 + 3

80

= 120

x

y

X

The

shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (iv). The feasible region is bounded.

Step III. The coordinates of the corner points O, A and C are (0, 0), (40, 0) and (0, 40) respectively.

Corner Point B: It is the point of intersection of boundary lines $2x + 3y = 120$ and $2x + y = 80$

Let us solve them for x and y. Subtracting $2y = 40 \Rightarrow y = 20$

Putting $y = 20$ in $2x + 3y = 120$; $2x + 60 = 120$
 $\Rightarrow 2x = 60 \Rightarrow x = 30$.

Therefore corner point B is (30, 20).

Step IV. Now, we evaluate Z at each corner point.

Corner Point Z = $7x + 10y$

$$\begin{aligned}
 &O(0, 0) \quad 0 \\
 &A(40, 0) \quad 280 \\
 &B(30, 20) \quad 410 = M \leftarrow \text{Maximum} \quad C(0, 40) \quad 400
 \end{aligned}$$

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By Corner Point Method, maximum $Z = 410$ at $(30, 20)$. Hence, the profit is maximum equal to ` 410 when 30 packages of screws A and 20 packages of screws B are produced in a day.

6. A cottage industry manufactures pedestal lamps and wooden shades, each requiring the use of a grinding/cutting machine and a sprayer. It takes 2 hours on grinding/cutting machine and 3 hours on the sprayer to manufacture a pedestal lamp. It takes 1 hour on the grinding/cutting machine and 2 hours on the sprayer to manufacture a shade. On any day, the sprayer is available for at the most 20 hours and the grinding/cutting machine for at the most 12 hours. The profit from the sale of a lamp is ` 5 and that from a shade is ` 3. Assuming that the manufacturer can sell all the lamps and shades that he produces, how should he schedule his daily production in order to maximise his profit?

Sol. Step I. Mathematical formulation of L.P.P.

Suppose the manufacturer produces x pedestal lamps and y wooden shades. The given data is condensed in the following table: Time on

Item	Number	grinding/	sprayer	Profit	cutting machine (hrs/item)
		(`/item)			
		(hrs/item)			

Pedestal lamps x 2 35 Wooden shades y 1 23 Total profit = $5x + 3y$

$$\text{Let } Z = 5x + 3y$$

We have the following mathematical model for the given problem:

Maximise $Z = 5x + 3y$... (i) subject to the constraints:

$$2x + y \leq 12 \quad (\text{Grinding/cutting machine constraint}) \quad \dots (ii)$$

[Given: Cutting/grinding machine is available for at the most (i.e., \leq) 12 hours]

$$3x + 2y \leq 20 \quad (\text{Sprayer constraint}) \quad \dots (iii) \quad [\text{Given: The sprayer}$$

is available for at the most 20 hours i.e., \leq

20] pedestal lamps and wooden shades can't be negative)

$$x, y \geq 0 \quad \dots (iv) \quad (\because \text{Number of}$$

Step II. The constraint (iv) $x, y \geq 0 \Rightarrow$ The feasible region is in first quadrant.

Table of values for the line $2x + y = 12$ of constraint (ii)

x	0	6
y	12	0

Let us draw the line joining the points $(0, 12)$ and $(6, 0)$. Let us test for origin $(x = 0, y = 0)$ in constraint (ii) $2x + y \leq 12$,

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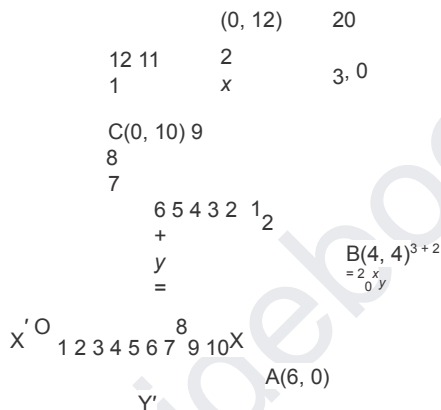
we have $0 \leq 12$ which is true.

\therefore Region for constraint (ii) is on the origin side of the line $2x + y = 12$.

Table of values for the line $3x + 2y = 20$ of constraint (iii)

Let us draw the line joining the points (0, 10) and

Y



Let us
test
for
origin
($x = 0$,
 $y = 0$)
in

constraint (iii) $3x + 2y \leq 20$, we have $0 \leq 20$ which is true. \therefore Region for constraint (iii) is on the origin side of the line $3x + 2y = 20$.

The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (iv). The feasible region is bounded.

Step III. The coordinates of the corner points O, A and C are (0, 0), (6, 0) and (0, 10) respectively.

Corner point B: It is the point of intersection of boundary lines $2x + y = 12$

and $3x + 2y = 20$

$2 \times$ First eqn. – Second eqn. gives

$$4x + 2y - 3x - 2y = 24 - 20 \Rightarrow x = 4.$$

Putting $x = 4$ in $2x + y = 12$, we have $8 + y = 12 \Rightarrow y = 4$.

\therefore Corner point B is (4, 4).

Step IV. Now, we evaluate Z at each corner point. . 27

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Corner Point $Z = 5x + 3y$

O(0, 0) 0

A(6, 0) 30

B(4, 4) 32 = M \leftarrow Maximum

C(0, 10) 30

By Corner Point Method, maximum $Z = 32$ at (4, 4).

Hence, the profit is maximum when 4 pedestal lamps and 4 wooden shades are manufactured. Maximum profit is ` 32. 7. A company manufactures two types of novelty souvenirs made of plywood.

Souvenirs of type A require 5 minutes each for cutting and 10 minutes each for assembling. Souvenirs of type B require 8 minutes each for cutting and 8 minutes each for assembling. There are 3 hours 20 minutes

available for cutting and 4 hours for assembling. The profit is ₹ 5 each for type A and ₹ 6 each for type B souvenirs. How many souvenirs of each type should the company manufacture in order to maximise the profit?
(Important)

Sol. Step I. Mathematical formulation of L.P.P.

Suppose the company manufactures x souvenirs of type A and y souvenirs of type B. The given data is condensed in the following table:

Time for cutting assembling (min/item) Time for Profit

Type Number (min/item) (min/item)

A x 5 10 5 B y 8 86 Total profit = $5x + 6y$

Let $Z = 5x + 6y$

We have the following mathematical model for the given problem:

Maximise $Z = 5x + 6y$... (i) subject to the constraints:

$5x + 8y \leq 200$ (Cutting constraint) ... (ii) [Given: (Maximum) time available for cutting is 3 hours, 20 minutes = $3 \times 60 + 20 = 200$ minutes]

$10x + 8y \leq 240$ (Assembling constraint) ... (iii) [Given: (Maximum) Time available for assembly is 4 hours = $4 \times 60 = 240$ minutes]

$x, y \geq 0$... (iv) (Number of souvenirs can't be negative)

Step II. Constraint (iv) $x, y \geq 0 \Rightarrow$ Feasible region is in first quadrant. Table of values for the line $5x + 8y = 200$ of constraint (ii)

x 0 40
 y 25 0

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Let
us

draw the line joining the points (0, 25) and (40, 0). Let us test for origin ($x = 0, y = 0$) in constraint (ii) $5x + 8y \leq 200$ we have $0 \leq 200$ which is true.

\therefore Region for constraint (ii) is on the origin side of the line $5x + 8y = 200$.

Table of values for the line $10x + 8y = 240$ of constraint (iii) x 0 24 y 30 0

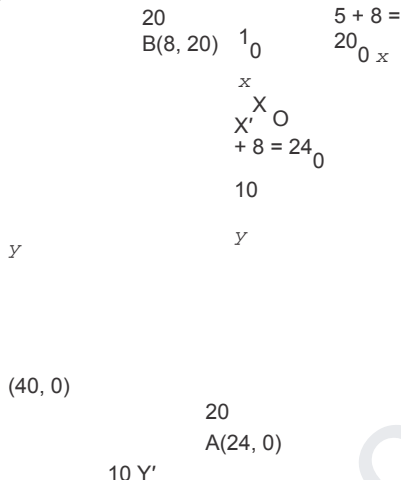
Let us draw the line joining the points (0, 30) and (24, 0). Let us test for origin ($x = 0, y = 0$) in constraint (iii) $10x + 8y \leq 240$, we have $0 \leq 240$ which is true.

\therefore Region for constraint (iii) is on the origin side of the line $10x + 8y = 240$.

Y

(0, 30)

C(0, 25)



The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (iv). The feasible region is bounded.

Step III. The coordinates of the corner points O, A and C are (0, 0), (24, 0) and (0, 25) respectively.

Corner point B: It is the point of intersection of the boundary lines

$$5x + 8y = 200 \text{ and } 10x + 8y = 240$$

$$\text{Subtracting, } -5x = -40 \Rightarrow x =$$

$$= 8.$$

Putting $x = 8$ in $5x + 8y = 200$, we have

$$40 + 8y = 200 \Rightarrow 8y = 160 \Rightarrow y =$$

$$= 20$$

\therefore Corner point B(8, 20).

Step IV. Now, we evaluate Z at each corner point. . 29 Class 12

$$O(0, 0) \quad 0$$

$$A(24, 0) \quad 120$$

$$B(8, 20) \quad 160 = M \leftarrow \text{Maximum}$$

$$C(0, 25) \quad 150$$

By Corner Point Method, maximum $Z = 160$ at $(8, 20)$. Hence, the profit is maximum when 8 souvenirs of type A and 20 souvenirs of type B are manufactured.

Maximum profit = ₹ 160.

8. A merchant plans to sell two types of personal computers a desktop model and a portable model that will cost ₹ 25,000 and ₹ 40,000

respectively. He estimates that the total monthly demand of computers will not exceed 250 units. Determine the number of units of each type of computers which the merchant should stock to get maximum profit if he does not want to invest more than ₹ 70 lakhs and if his profit on the desktop model is ₹ 4500 and on portable model is ₹ 5000.

Sol. Step I. Mathematical Formulation of L.P.P.

Suppose the merchant stocks x units of desktop model and y units of portable model. The given data is condensed in the following table.

Type	Number	Cost (₹/unit)	Profit (₹/unit)
Desktop	x	25000	4500
Portable	y	40000	5000
Total profit		$= 4500x + 5000y$	

$$\text{Let } Z = 4500x + 5000y$$

We have the following mathematical model for the given problem:

Maximise profit $Z = 4500x + 5000y$... (i) subject to the constraints:

$x + y \leq 250$ (Demand constraint) ... (ii) [Given: Total monthly demand of computers will not exceed 250 i.e., ≤ 250]

$$25000x + 40000y \leq 70,00,000$$

[Given: He does not want to invest more than ₹ 70 lakhs = ₹ 70

$$\times 100,000]$$

Dividing every term by 5000,

$$\text{or } 5x + 8y \leq 1400 \text{ (Investment constraint) ... (iii)}$$

$$x, y \geq 0 \text{ ... (iv)}$$

(∵ Number of computers can't be negative)

Step II. Constraint (iv) $x, y \geq 0 \Rightarrow$ Feasible region is in first

quadrant. . 30

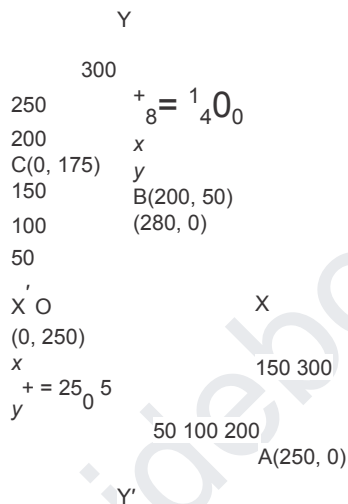
Let us draw the line joining the points (0, 250) and (250, 0). Let us test for origin ($x = 0, y = 0$) in constraint (ii) $x + y \leq 250$, we have $0 \leq 250$ which is true.

\therefore Region for constraint (ii) is on the origin side of the line $x + y = 250$.

Table of values for the line $5x + 8y = 1400$ of constraint (iii) $x \leq 280$
 $y \leq 175$

Let us draw the line joining the points (0, 175) and (280, 0). Let us test for origin (0, 0) in constraint (iii), $5x + 8y \leq 1400$, we have $0 \leq 1400$ which is true.

\therefore Region for constraint (iii) is on the origin side of the line $5x + 8y = 1400$.



The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (iv). The feasible region is bounded.

Step III. The coordinates of the corner points O, A and C are (0, 0), (250, 0) and (0, 175) respectively.

Corner point B: It is the point of intersection of boundary lines: $x + y = 250$ and $5x + 8y = 1400$

Second Eqn. $- 5 \times$ 1st equation gives

$$5x + 8y - 5x - 5y = 1400 - 1250$$

$$\text{or } 3y = 150 \Rightarrow \text{Putting } y = 50 \text{ in } x + y = 250,$$

$$y = 50$$

$$\text{we have } x + 50 = 250 \Rightarrow x = 200$$

\therefore Corner point B is (200, 50).

Step IV. Now, we evaluate Z at each corner point.

$$\text{Corner Point } Z = 4500x + 5000y$$

$$O(0, 0) \quad 0$$

$$A(250, 0) \quad 11,25,000$$

$$B(200, 50) \quad 11,50,000 = M \leftarrow \text{Maximum } C(0, 175) \quad 8,75,000$$

By Corner Point Method, maximum $Z = 11,50,000$ at $(200, 50)$. Hence, the merchant should stock 200 units of desktop model and 50 units of portable model for a maximum profit of ₹ 11,50,000.

9. A diet is to contain at least 80 units of vitamin A and 100 units of minerals. Two foods F_1 and F_2 are available. Food F_1 costs ₹ 4 per unit and food F_2 costs ₹ 6 per unit. One unit of food F_1 contains 3 units of vitamin A and 4 units of minerals. One unit of food F_2 contains 6 units of vitamin A and 3 units of minerals. Formulate this as a linear programming problem. Find the minimum cost for diet that consists of mixture of these two foods and also meets the minimal nutritional requirements.

Sol. Step I. Mathematical formulation of L.P.P.

Suppose the diet contains x units of food F_1 and y units of food F_2 . The given data is condensed in the following table:

Type	Number	Cost
Vitamin A	Minerals of Food of units (unit)	(units) (units)
F_1	3	4
F_2	6	3
Total cost	$= 4x + 6y$	

$$\text{Let } Z = 4x + 6y$$

We have the following mathematical model for the given problem.

Minimise $Z = 4x + 6y$... (i) subject to the constraints:

$3x + 6y \geq 80$ (Vitamin A constraint) ... (ii) [Given: At least i.e., \geq 80 units of vitamin A]

$4x + 3y \geq 100$ (Mineral constraint) ... (iii) [Given: At least i.e., \geq 100 units of minerals]

$$x, y \geq 0$$

(∵ Units of vitamins and minerals can't be negative) ... (iv) Step

II. The constraint (iv) $x, y \geq 0$.

⇒ Feasible region is in first quadrant.

Table of values for the line $3x + 6y = 80$ of constraint (ii)

x	y
0	
	0

0

y

Let us draw the line joining the points

$$(0, 0) \text{ and } (25, 0)$$

Let us test for origin ($x = 0, y = 0$) in constraint (ii) $3x + 6y \geq 80$, we have $0 \geq 80$ which is not true.

\therefore Region for constraint (ii) is the half-plane not containing the origin i.e., region on the non-origin side of the line $3x + 6y =$

80. Table of values for the line $4x + 3y = 100$ of constraint (iii)

$$x \ 0 \ 25$$

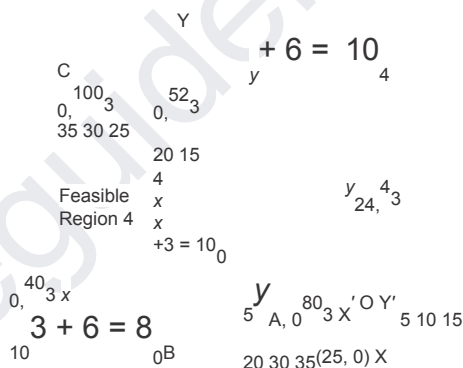
$$y \ 0$$

Let us draw the line joining the points

$$(0, 0) \text{ and } (25, 0).$$

Let us test for origin ($x = 0, y = 0$) in constraint (iii) $4x + 3y \geq 100$, we have $0 \geq 100$ which is not true.

\therefore Region for constraint (iii) is the half-plane again on the non origin side of the line $4x + 3y = 100$.



The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (iv).

The feasible region is unbounded.

Step III. The coordinates of the corner points A and C are

and

respectively.

$$\begin{aligned} & \frac{1}{2} \times \frac{1}{2} \\ & \frac{1}{2} \times \frac{1}{2} \end{aligned}$$

To find corner point B: Corner point B is the point of intersection of the boundary lines

$$3x + 6y = 80 \text{ and } 4x + 3y = 100$$

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First Eqn. $- 2 \times$ Second eqn. gives

$$3x + 6y - 8x - 6y = 80 - 200$$

$$\text{or } -5x = -120 \Rightarrow x =$$

-

$$= 24$$

Putting $x = 24$ in $3x + 6y = 80$, we have

$$72 + 6y = 80 \Rightarrow 6y = 8 \Rightarrow y = \frac{4}{3} \therefore \text{Corner}$$

point B is

$$\begin{aligned} & \frac{1}{2} \times \frac{1}{2} \\ & \frac{1}{2} \times \frac{1}{2} \end{aligned}$$

Step IV. Now, we evaluate Z at each corner point.

Corner Point Z = $4x + 6y$

A

$$\begin{aligned} & \frac{1}{2} \times \frac{1}{2} \\ & \frac{1}{2} \times \frac{1}{2} \end{aligned}$$

B

$$\begin{aligned} & \frac{1}{2} \times \frac{1}{2} \\ & \frac{1}{2} \times \frac{1}{2} \end{aligned} 104 = m \leftarrow \text{Smallest}$$

C

$$\begin{aligned} & \frac{1}{2} \times \frac{1}{2} \\ & \frac{1}{2} \times \frac{1}{2} \end{aligned} 200$$

From this table, we find that 104 is the smallest value of Z at the

$$\frac{1}{2} \times \frac{1}{2}$$

corner B

$$\frac{1}{2} \times \frac{1}{2}$$

Step V. Since the feasible region is unbounded, 104 may or may not

be the minimum value of Z . To decide this, we graph the inequality $Z < m$ i.e., $4x + 6y < 104$.

Table of values for the line $4x + 6y = 104$ (of constraint $Z < m$ i.e., $4x + 6y < 104$)

x 0 26

0

y

Let us draw the dotted line joining the points

$(0, 0)$ and $(26, 0)$.

$(26, 0)$ not being marked in the graph because it is very close to the

point $(26.7, 0)$ already marked and $(26, 0)$ is slightly to the

left of

$(26.7, 0)$

The line is shown dotted because equality sign is absent in the constraint $Z < m$.

Let us test for origin $(x = 0, y = 0)$ in constraint $Z < m$ i.e., $4x + 6y < 104$, we have $0 < 104$ which is true.

\therefore Region for constraint $Z < m$ i.e., $4x + 6y < 104$ is the origin side of the line $4x + 6y = 104$

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We observe that the half plane determined by $Z < m$ has no point in common with the feasible region. Hence $m = 104$ is the

minimum value of Z attained at the point B

$(24, 10)$

\therefore Minimum cost is ` 104 when 24 units of food F_1 are mixed with units of food F_2 .

10. There are two types of fertilisers F_1 and F_2 . F_1 consists of 10% nitrogen and 6% phosphoric acid and F_2 consists of 5% nitrogen and 10% phosphoric acid. After testing the soil conditions, a farmer finds that she needs atleast 14 kg of nitrogen and 14 kg of phosphoric acid for her crop. If F_1 costs ` 6/kg and F_2 costs ` 5/kg, determine how much of each type of fertiliser should be used so that nutrient requirements are met at a minimum cost. What is the minimum cost?

Sol. Step I. Mathematical formulation of L.P.P.

Suppose the farmer uses x kg of fertiliser F_1 and y kg of fertiliser F_2 . The given data is condensed in the following table.

Fertiliser	Quantity	Nitrogen content (%)	Phosphoric content (%)	Cost (₹/kg)
F_1	10%	6%	6%	6
F_2	5%	10%	5%	5

$$F_1 \times 10\% \text{ Nitrogen} + F_2 \times 5\% \text{ Nitrogen} \geq 14\% \text{ Nitrogen}$$

$$\text{Total cost} = 6x + 5y$$

$$\text{Let } Z = 6x + 5y$$

We have the following mathematical model for the given problem:

Minimise $Z = 6x + 5y$... (i) subject to the constraints:

$$x + y \geq 14$$

$$y \geq 14$$

[Given: She needs at least i.e., ≥ 14 kg of nitrogen for her crops]

Multiplying by 100 and dividing by 5,

$$2x + y \geq 280 \text{ (Nitrogen constraint) ... (ii)}$$

$$y \geq 14$$

[Given: She needs at least 14 kg of phosphoric acid for her crops]

Multiplying by 100 and dividing by 2,

$$3x + 5y \geq 700 \text{ (Phosphoric acid constraint) ... (iii)}$$

$$x, y \geq 0 \text{ ... (iv)}$$

Quantity of Nitrogen and Phosphoric acid can't be negative) Step II.

Constraint (iv) $x, y \geq 0$.

⇒ Feasible region is in first quadrant.

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Table of values for the line $2x + y = 280$ of constraint (ii)

$$x \ 0 \ 140$$

$$y \ 280 \ 0$$

Let us draw the line joining the points (0, 280) and (140, 0). Let us test for origin ($x = 0, y = 0$) in constraint (ii) $2x + y \geq 280$, we have $0 \geq 280$ which is not true.

∴ Region for constraint (ii) is the half-plane not containing the origin i.e., region on the non-origin side of the line $2x + y = 280$. Table of

values for the line $3x + 5y = 700$ corresponding to constraint (iii)

$x = 0$

$y = 140$

Let us draw the line joining the points $(0, 140)$ and $(280, 0)$.

Let us test for origin $(x = 0, y = 0)$ in constraint (iii) $3x + 5y \geq 700$, we have $0 \geq 700$ which is not true.

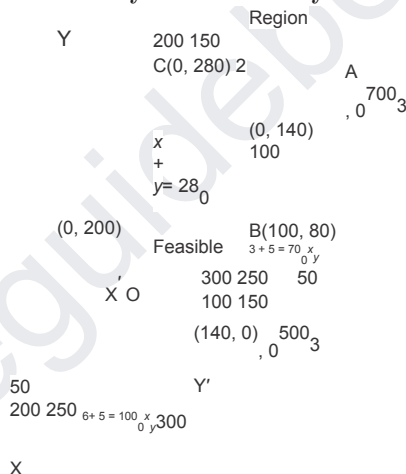
\therefore Region for constraint (iii) is again on the non-origin side of the line $3x + 5y = 700$.

The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (iv). The feasible region is unbounded.

Step III. The coordinates of the corner points. A and C are

$(0, 140)$ and $(0, 280)$ respectively.

To find corner point B: Let us solve the equations of bounding lines $2x + y = 280$ and $3x + 5y = 700$ for x and y .



Second eqn. $- 5 \times$ first eqn. gives

$$3x + 5y - 10x - 5y = 700 - 1400$$

$$\Rightarrow -7x = -700 \Rightarrow x =$$

—

$$= 100$$

Putting $x = 100$ in $2x + y = 280$, we have

$200 + y = 280 \Rightarrow y = 80 \therefore$ Corner point B is (100, 80). Step IV. Now, we evaluate Z at each corner point.

$$\text{Corner Point } Z = 6x + 5y$$

$$\begin{array}{l} A \\ \square \square \square \\ \square \square 1400 \end{array}$$

$$B(100, 80) \quad 1000 = m \leftarrow \text{Smallest}$$

$$C(0, 280) \quad 1400$$

From this table, we find that 1000 is the smallest value of Z at the corner B(100, 80). Since the feasible region is unbounded, 1000 may or may not be the minimum value of Z.

Step V. To decide this, we graph the inequality $Z < m$ i.e., $6x + 5y < 1000$.

Table of values for the line $6x + 5y = 1000$ (for constraint $Z < m$ i.e., $6x + 5y < 1000$)

$$x \quad 0$$

$$y \quad 200 \quad 0$$

Let us draw the dotted line joining the points (0, 200) and \square
 $\square \square \square \square$.

The line is drawn dotted because equality sign is absent in the constraint $Z < m$.

We observe that the half-plane determined by $Z < m$ has no point in common with the feasible region. Hence, $m = 1000$ is the minimum value of Z attained at the point B(100, 80). \therefore Minimum cost is ` 1000 when the farmer uses 100 kg of fertiliser F_1 and 80 kg of fertiliser F_2 .

11. The corner points of the feasible region determined by the following system of linear inequalities:

$2x + y \leq 10$, $x + 3y \leq 15$, $x, y \geq 0$ are (0, 0), (5, 0), (3, 4) and (0, 5). Let $Z = px + qy$, where $p, q > 0$. Condition on p and q so that the maximum of Z occurs at both (3, 4) and (0, 5) is (A) $p = q$ (B) $p = 2q$ (C) $p = 3q$ (D) $q = 3p$.

Sol. We evaluate Z at each corner point.

$$\begin{array}{l} \text{Corner Point } Z = px + qy, \\ p > 0, q > 0 \end{array}$$

$$(0, 0) \quad 0$$

$$(5, 0) \quad 5p$$

$$(3, 4) \quad 3p + 4q$$

$$(0, 5) \text{ } 5q = M \leftarrow \text{Maximum}$$

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\therefore Maximum of Z occurs at both (3, 4) and (0, 5) (given) $\therefore 3p +$

$$4q = 5q$$

$$\therefore q = 3p$$

Hence, the correct option is (D).

MISCELLANEOUS EXERCISE

1. (Refer to Example 9, NCERT Page 521). How many packets of each food should be used to maximise the amount of vitamin A in the diet? What is the maximum amount of vitamin A in the diet?

Sol. (NCERT Page 521), we find that Z is maximum at the point (40, 15). Hence, the amount of vitamin A under the constraints given in the problem will be maximum if 40 packets of food P and 15 packets of food Q are used in the special diet.

The maximum amount of vitamin A will be 285 units.

2. A farmer mixes two brands P and Q of cattle feed. Brand P costing ₹ 250 per bag, contains 3 units of nutritional element A, 2.5 units of element B and 2 units of element C. Brand Q costing ₹ 200 per bag contains 1.5 units of nutritional element A, 11.25 units of element B, and 3 units of element C. The minimum requirements of nutrients A, B and C are 18 units, 45 units and 24 units respectively. Determine the number of bags of each brand which should be mixed in order to produce a mixture having a minimum cost per bag? What is the minimum cost of the mixture per bag?

Sol. Step I. Mathematical Formulation of L.P.P.

Suppose the farmer mixes x bags of brand P and y bags of brand Q.

The given data is condensed in the following table.

Brand	Number	Cost	Element A	Element B	Element C	of bags (₹/bag)
(units/bag)	(units/bag)	(units/bag)	P	x	250	3
			2.5	2	Q	y
			200	1.5	11.25	3
Total						
cost = 250x + 200y						

$$\text{Let } Z = 250x + 200y$$

We have the following mathematical model for the given problem:

Minimise $Z = 250x + 200y$... (i) subject to the constraints:

$$3x + 1.5y \geq 18$$

[Given: Minimum requirement of nutritional element A is 18 units i.e., ≥ 18 units]

$$\text{or } 3x + y \geq 18$$

Multiplying by 10 and dividing by 15,

$$\text{or } 2x + y \geq 12 \text{ (Nutritional element A constraint)...(ii) } 2.5x + 11.25y \geq 45$$

[Given: Minimum requirement of nutritional element B is, 45

units . 39

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i.e., ≥ 45 units]

$$\text{or } x +$$

$$y \geq 45$$

Multiplying by 100 and dividing by 125,

$$\text{or } 2x + 9y \geq 36 \text{ ... (iii) (Nutritional element B constraint)}$$

$$2x + 3y \geq 24 \text{ (Nutritional element C constraint) ... (iv) [Given:}$$

Minimum requirement of nutritional element C is 24 units i.e., ≥ 24 units]

$$x, y \geq 0 \text{ (Number of bags can't be negative)...(v)}$$

Step II. Constraint (v) $x, y \geq 0$

\Rightarrow Feasible region is in first quadrant.

Table of values for the line $2x + y = 12$ of constraint (ii)

$x \quad 0 \quad 6$
 $y \quad 12 \quad 0$

Draw the straight line joining the points (0, 12) and (6, 0). Let us test for origin ($x = 0, y = 0$) in constraint $2x + y \geq 12$, we have $0 \geq 12$ which is not true.

\therefore Region for constraint (ii) $2x + y \geq 12$ is the half-plane not containing the origin i.e., region on the non-origin side of the line $2x + y = 12$. Table of values for the line $2x + 9y = 36$ for constraint (iii)

$x \quad 0 \quad 18$
 $y \quad 4 \quad 0$

Let us draw the line joining the points (0, 4) and (18, 0). Let us test for origin ($x = 0, y = 0$) in constraint (iii) $2x + 9y \geq 36$, we have $0 \geq 36$ which is not true.

\therefore Region for constraint (iii) is the region on the non-origin side of the line $2x + 9y = 36$.

Table of values for the line $2x + 3y = 24$ for constraint (iv)

$x \quad 0 \quad 12$
 $y \quad 8 \quad 0$

Draw the line joining the points (0, 8) and (12, 0). Let us test for origin ($x = 0, y = 0$) in constraint (iii) $2x + 3y \geq 24$, we have $0 \geq 24$ which is not true.

\therefore Region for constraint (iii) $2x + 3y \geq 24$ is again the region on the non-origin side of the line $2x + 3y = 24$.

The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (v). The feasible region is unbounded.

Step III. The coordinates of the corner points A and D are (18, 0) and (0, 12) respectively.

Corner point B: It is the point of intersection of the lines $2x + 3y = 24$ and $2x + 9y = 36$

Subtracting $-6y = -12 \Rightarrow y =$

$-$

$- = 2$

$.40$

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Y

5

x

$+4 = 3_9$

y

14 12 10
8
(0, 8) 6
4
(0, 4) 2
D(0, 12) Feasible
Region

$$2x + 9y = 36$$

C(3, 6)

B(9, 2)

X' O 2 4 6 8 10 12 14 16 X

X

A(18, 0)

x
y

Y'
(6, 0) (12, 0) +
y
= 12

$$2x + 3y = 24$$

Putting $y = 2$ in $2x$

+ $3y = 24$, we have
 $2x + 6 = 24 \Rightarrow 2x = 18 \Rightarrow x = 9$
 \therefore Corner point B
is (9, 2).

Subtracting $-2y = -12 \Rightarrow y =$

Putting $y = 6$ in $2x + y = 12$, we have

\therefore Corner point C is (3, 6).

Corner Point Z = 250x + 200y

B(9, 2) 2650

From this table, we find that 1950 is the smallest value of Z at the corner C(3, 6). Since the feasible region is unbounded, 1950 may or may not be the minimum value of Z.

$$250x + 200y < 1950 \text{ or } 5x + 4y < 39.$$

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$$x = 7.8$$

$$y = 9.75$$

Let us draw the dotted line joining the points (0, 9.75) and (7.8, 0). The line is to be shown dotted because equality sign is absent in the constraint $Z < m$ i.e., in $5x + 4y < 39$.

Let us test for origin ($x = 0, y = 0$) in this constraint, we have $0 < 39$ which is true.

\therefore Region for constraint $Z < m$ i.e., $5x + 4y < 39$ is towards the origin side of the line.

We observe that the half plane determined by $Z < m$ has no point in common with the feasible region. Hence $m = 1950$ is the minimum value of Z attained at the point C(3, 6).

\therefore Minimum cost is ` 1950 when 3 bags of brand P and 6 bags of brand Q are mixed.

3. A dietician wishes to mix together two kinds of food X and Y in such a way that the mixture contains at least 10 units of vitamin A, 12 units of vitamin B and 8 units of vitamin C. The vitamin contents of one kg food is given below:

Food	Vitamin A	Vitamin B	Vitamin C	X	Y
	10	12	8	2	3
	2	1			

One kg of food X costs ` 16 and one kg of food Y costs ` 20. Find the least cost of the mixture which will produce the required diet? Sol. Step I. Mathematical Formulation of L.P.P.

Let the dietician mix x kg of food X and y kg of food Y. The given data is condensed in the following table.

Food	Quantity	Vitamin A	Vitamin B	Vitamin C	Cost (kg)	(units/kg)
	(units/kg)	(units/kg)	(units/kg)	(units/kg)	X	Y
	10	12	8	16	2	3
	2	1		20	2	2
	120					
	Total cost				$16x + 20y$	

Let $Z = 16x + 20y$

We have the following mathematical model for the given problem:

Minimise $Z = 16x + 20y$... (i) subject to the constraints:

$x + 2y \geq 10$ (Vitamin A constraint) ... (ii) [Given: The mixture

contains at least 10 units (i.e., ≥ 10) of vitamin A] $2x + 2y \geq 12$

[Given: The mixture contains at least 12 units (i.e., ≥ 12) of vitamin

B] or $x + y \geq 6$ (Vitamin B constraint) ... (iii) $3x + y \geq 8$ (Vitamin C

constraint) ... (iv) [Given: The mixture contains at least 8 units (i.e.,

≥ 8) of vitamin C]

$x,$

$y \geq$

\therefore Quantities of food can't be negative) ... (v) The constraint (v), $x, y \geq 0 \Rightarrow$ Feasible region is in first quadrant.

Table of values for the line $x + 2y = 10$ of constraint (ii). $x \geq 0, y \geq 0$

Let us draw the line joining the points (0, 5) and (10, 0). Let us test for origin ($x = 0, y = 0$) in constraint (ii) $x + 2y \geq 10$, we have $0 \geq 10$ which is not true.

\therefore Region for constraint (ii) is the half-plane not containing the origin i.e., region on the non-origin side of the line $x + 2y = 10$. Table of values for the line $x + y = 6$ of constraint (iii). $x \geq 0, y \geq 0$

$y \geq 0$

Let us draw the line joining the points (0, 6) and (6, 0). Let us test for origin ($x = 0, y = 0$) in constraint $x + y \geq 6$, we have $0 \geq 6$ which is not true.

\therefore Region for constraint (iii) is the half-plane not containing the origin i.e., region on the non-origin side of the line $x + y = 6$. Table of values

for the line $3x + y = 8$ of constraint (iv). $x \geq 0$

$y \geq 0$

Let us draw the line joining the points (0, 8) and $\square \square \square \square$.

Let us test for origin ($x = 0, y = 0$) in constraint (iv) $3x + y \geq 8$, we have $0 \geq 8$ which is not true.

\therefore Region for constraint (iv) also is on the non-origin side of the line $3x + y = 8$.

Y

D(0, 8) Feasible

(0, 5)

4 3

3

x

= 8

Region

$\frac{1}{0}$

(0, 6)

C(1, 5)

$x + y + 2 =$

y

2 1

8

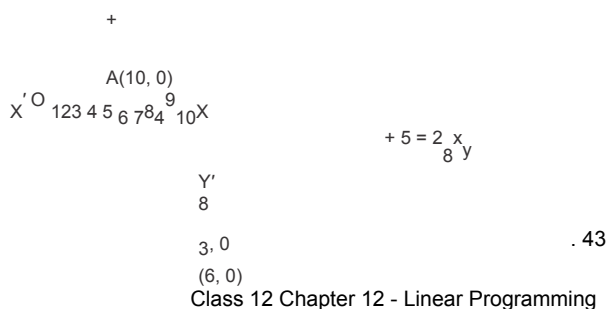
7

6

y

+B(2, 4) x

= 6



The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (v). The feasible region is unbounded.

Step III. The coordinates of the corner points A and D are (10, 0) and (0, 8) respectively.

Corner point B: It is the point of intersection of bounding lines $x + 2y = 10$ and $x + y = 6$

Subtracting $y = 4$

Putting $y = 4$ in $x + 2y = 10$, $x + 8 = 10 \Rightarrow x = 2 \therefore$ Corner point B is (2, 4).

Corner point C: It is the point of intersection of bounding lines $x + y = 6$ and $3x + y = 8$

Subtracting $-2x = -2$ or $x =$

$$- \\ _ = 1$$

Putting $x = 1$ in $x + y = 6$, $1 + y = 6 \Rightarrow y = 5$

\therefore Corner point C is (1, 5).

Step IV. Now, we evaluate Z at each corner point.

Corner Point $Z = 16x + 20y$

A(10, 0) 160

B(2, 4) 112 = m \leftarrow Smallest C(1, 5) 116

D(0, 8) 160

From this table, we find that 112 is the smallest value of Z at the corner B(2, 4). Since the feasible region is unbounded, 112 may or may not be the minimum value of Z.

Step V. To decide this, we graph the inequality $Z < m$ i.e., $16x + 20y < 112$ or $4x + 5y < 28$.

Table of values for the line $4x + 5y = 28$ (of constraint $Z < m$ i.e., $4x + 5y < 28$).

$$x \ 0 \ 7$$

$$= 5.6 \ 0$$

$$y$$

Let us draw the dotted line joining the points (0, 5.6) and (7, 0). The line is drawn dotted because equality sign is absent in the constraint $Z < m$ i.e., $4x + 5y < 28$.

Let us test for origin ($x = 0, y = 0$) in constraint $4x + 5y < 28$, we have $0 < 28$ which is true.

\therefore Region for constraint $Z < m$ i.e., $4x + 5y < 28$ is on the origin side of the line $4x + 5y = 28$.

We observe that the half-plane determined by $Z < m$ has no point in common with the feasible region. Hence, $m = 112$ is the minimum value of Z attained at the point B(2, 4). \therefore Minimum cost of the mixture is ` 112 when 2 kg of food X and 4 kg of food Y are mixed.

4. A manufacturer makes two types of toys A and B. Three machines are needed for this purpose and the time (in minutes) required for each toy on the machines is given below:

Types of Toys \ Machines

	A	12	18	6
	B	6	9	5

I II III

Each machine is available for a maximum of 6 hours per day. If the profit on each toy of type A is ` 7.50 and that on each toy of type B is ` 5, show that 15 toys of type A and 30 of type B should be manufactured in a day to get maximum profit.

Sol. Step I. Mathematical formulation of L.P.P.

Let the manufacturer make x toys of type A and y toys of type B. The given data is condensed in the following table.

Types of	Number	Time (min/toy)	on machines	Profit	toy of	toys	I	II	III
(/toy)	A	12	18	6	7.50	B	6	9	5
Total profit =					$7.50x + 5y$				

Let $Z = 7.50x + 5y$

We have the following mathematical model for the given problem:

Maximise $Z = 7.50x + 5y$... (i) subject to the constraints:

$$12x + 6y \leq 360$$

[Given: Each of machines I, II, III is available for a maximum of 6 hours = $6 \times 60 = 360$ minutes]

or $2x + y \leq 60$ (Machine I constraint) ... (ii) $18x + 0y \leq 360$

or $x \leq 20$ (Machine II constraint) ... (iii) $6x + 9y \leq 360$

or $2x + 3y \leq 120$ (Machine III constraint) ... (iv) $x, y \geq 0$... (v)

(\therefore Number of toys can't be negative)

Step II. Constraint (v) $x, y \geq 0$.

\Rightarrow Feasible region is in first quadrant.

Table of values for the line $2x + y = 60$ of constraint (ii). x 0 30

y 60 0

Let us draw the line joining the points (0, 60) and (30, 0). Let us test for

origin ($x = 0, y = 0$) in constraint (ii) $2x + y \leq 60$, we have $0 \leq 60$ which is true. Therefore region for constraint (ii) is on the origin side of the line $2x + y = 60$.

Region for constraint (iii) $x \leq 20$

We know that graph of the line $x = 20$ is a vertical line (parallel to y -axis) at a distance of 20 units along OX.

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∴

Region for $x \leq 20$ is the region on the left side of the line $x = 20$.

Table of values for the line $2x + 3y = 120$ of constraint (iv). x 0 60 y 40 0

Let us draw the line joining the points (0, 40) and (60, 0). Let us test for origin ($x = 0, y = 0$) in constraint (iv) $2x + 3y \leq 120$, we have $0 \leq 120$ which is true.

∴ Region for constraint (iv) is on the origin side of the line $2x + 3y = 120$.

The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (v). The feasible region is bounded. **Step III.** The coordinates of the corner points O, A and D are (0, 0), (20, 0) and (0, 40) respectively.

Corner point B: It is the point of intersection of bounding lines $2x + y = 60$ and $x = 20$

Putting $x = 20$ in $2x + y = 60$, we have $40 + y = 60$ or $y = 20$. ∴ Corner point B is (20, 20).

Corner point C: It is the point of intersection of bounding lines $2x + y = 60$ and $2x + 3y = 120$

Subtracting $-2y = -60$ or $y =$

—

— = 30

Putting $y = 30$ in $2x + y = 60$, we have

$2x + 30 = 60 \Rightarrow 2x = 30 \Rightarrow x = 15$.

∴ Corner point C is (15, 30).

Step IV. Now, we evaluate Z at each corner point.

Corner Point $Z = 7.50x + 5y$

O(0, 0) 0

A(20, 0) 150

B(20, 20) 250

C(15, 30) 262.50 = M ← Maximum D(0, 40) 200

By Corner Point Method, maximum $Z = 262.50$ at (15, 30). √

$(0, 60)$ 50 $D(0, 40)$ 30 20
 10
 $x = 20$
 $C(15, 30)$ $B(20, 20)$
 $(20, 0)$ $(30, 0)$ $(60, 0)$

$X \leq 10$

$$+40 \cdot 50 \cdot 2 \cdot x + 3 = 12 \cdot y$$

x

A

2

$$y = 60$$

Y'

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\therefore For maximum profit, 15 toys of type A and 30 toys of type B should be manufactured.

5. An aeroplane can carry a maximum of 200 passengers. A profit of ₹ 1000 is made on each executive class ticket and a profit of ₹ 600 is made on each economy class ticket. The airline reserves at least 20 seats for executive class. However, at least 4 times as many passengers prefer to travel by economy class than by the executive class. Determine how many tickets of each type must be sold in order to maximise the profit for the airline. What is the maximum profit?

Sol. Step I. Let us formulate the L.P.P. mathematically. Let the number of executive class tickets sold be x and the number of economy class tickets sold be y .

The aeroplane can carry a maximum of 200 passengers. $\Rightarrow x + y \leq 200$

$$x + y \leq 200$$

At least 20 seats are reserved for executive class $\Rightarrow x \geq 20$ Number of passengers in economy class is at least 4 times the number of passengers in executive class.

$$\Rightarrow y \geq 4x$$

Profit from x executive class tickets at the rate of ₹ 1000 per ticket = ₹ 1000 x

Profit from y economy class tickets at the rate of ₹ 600 per ticket = ₹ 600 y .

Let the total profit (in ₹) be denoted by P , then $P = 1000x + 600y \therefore$

Also $x \geq 0$ and $y \geq 0$ [Number of tickets can't be negative.] **Step II.** The reader is suggested to draw the graphs of constraints $x + y \leq 200$ and $x \geq 20$ for himself or herself and compare them with

the constraint $y \geq 4x$. The line $y = 4x$ is the line $y = 20$ when $x = 5$.

Put $x = 20$, $y = 80 \therefore$ Point is $(20, 80)$.
The graph of line $y = 4x$ is through the origin $(0, 0)$ and point $(20, 80)$.

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Test for the point $(1, 0)$.

Putting $x = 20$, $y = 4 \times 20 = 80$

\therefore Vertex A is (20, 80)

Corner (or vertex) B is the point of intersection of the lines $y = 4x$ and $x + y = 200$.

Putting $y = 4x$, $x + 4x = 200$ or $5x = 200$

$$\therefore x = 40 \text{ and therefore } y = 4x = 4(40) = 160$$

\therefore Vertex B is (40, 160)

Corner (or vertex) C is the point of intersection of the lines $x = 20$ and $x + y = 200$.

Putting $x = 20$, $20 + y = 200$

$$\therefore y = 180$$

\therefore Vertex C is (20, 180)

Step IV. Objective function is $P = 1000x + 600y$.

At A (20, 80); $P = 1000(20) + 600(80) = 68000$

At B (40, 160); $P = 1000(40) + 600(160) = 136000$

At C (20, 180); $P = 1000(20) + 600(180) = 128000$

We see that P is maximum at B where $x = 40$, $y = 160$. \therefore The airline should sell 40 executive class tickets and 160 economy class tickets to maximise profit.

Also, maximum profit = The value of P at B = ` 136000. 6. Two godowns A and B have grain capacity of 100 quintals and 50 quintals respectively. They supply to 3 ration shops, D, E and F whose requirements are 60, 50 and 40 quintals respectively. The cost of transportation per quintal from the godowns to the shops are given in the following table:

Transportation Cost per quintal (in `)

From / To A B

D	6	4
E	3	2
F	2.50	3

How should the supplies be transported in order that the transportation cost is minimum? What is the minimum cost? Sol. Step I. Mathematical formulation of L.P.P.

Let x quintals and y quintals of grain be transported from . 48

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godowns A to ration shops D and E respectively. Then $100 - (x + y)$ quintals will be transported to ration shop F.

Clearly, $x \geq 0$, $y \geq 0$ and $100 - x - y \geq 0$ ($\Rightarrow 100 \geq x + y$). \therefore Amounts (in Quintals) of grain can't be negative) i.e., $x \geq 0$, $y \geq 0$ and $x + y \leq 100$

Now, the requirement of x quintals are transported from godown A, the requirement of x quintals are transported from godown A 100 - x

remaining $(60 - x)$ quintals need to be transported from godown B. 60 50

$\therefore 3, 6, 2.50$

y
 y

Similarly, $(50 - y)$ and $40 - (100 - x - y)$

Shop

$$= x + y - 60 \text{ quintals}$$

F
D

E
need to be transported from godown B to shops E and F respectively.
60 - x B y + - 60 x

i.e., $x \leq 60$, $y \leq 50$ and $x + y \geq 60$ Total transportation cost Z is

Clearly, $60 - x \geq 0$, $50 - y \geq 0$ (i.e., $60 \geq x$, $50 \geq y$) and $x + y$ given by Godown

$$Z = 6x + 3y + (100 - x - y) + 4(60 - x) + 2(50 - y) + 3(x + y - 60)$$

$$x + y + 410 = (5x + 3y + 820) =$$

We have the following mathematical model for the given problem:
Minimise $Z = (5x + 3y + 820)$... (i) subject to the constraints:

$x \geq 0$, $y \geq 0$... (ii) $x + y \leq 100$... (iii) $x \leq 60$... (iv) $y \leq 50$... (v) $x + y \geq 60$... (vi) Step II. Constraint (ii) $x \geq 0$, $y \geq 0$.
⇒ Feasible region is in first quadrant.

Table of values for the line $x + y = 100$ of constraint (iii). $x \geq 0$ $y \geq 0$

y 100 0

Let us draw the straight line joining the points (0, 100) and (100, 0). Let us test for origin ($x = 0$, $y = 0$) in constraint (ii) $x + y \leq 100$, we have $0 \leq 100$ which is true.

\therefore Region for constraint (ii) is on the origin side of the line $x + y = 100$. Region for constraint (iv) $x \leq 60$

We know that graph of the line $x = 60$ is a vertical line (parallel to y-axis) at a distance of 60 units along OX.

\therefore Region for constraint $x \leq 60$ is the region on the left side of the line $x = 60$.

Region for constraint (v) $y \leq 50$

We know that graph of the line $y = 50$ is a horizontal line (parallel to x-axis) at a distance of 50 units along OY. \therefore Region for constraint $y \leq 50$ is below the line $y = 50$. Finally, Table of values for the line $x + y = 60$ of constraint (vi). $x \geq 0$

$y \geq 0$

Let us draw the line

$x + y = 100$

joining the points $(0, 0)$ and $(60, 0)$ and $(0, 100)$. Let us test for origin ($x = 0, y = 0$) in constraint (vi) $x + y \geq 60$. we have $0 \geq 60$

$(0, 100)$

$x = 60$

which is not true. \therefore
S(10, 50) $y = 50$ R(50, 50)

Region for

50

origin i.e., $(0, 0)$

$= 60$

x

$+$

y

Q(60, 40)

constraint (vi) is the half plane not containing the

region on the

$x + y = 100$ shaded

the $x'O'Y'$ figure is

$(100, 0)$

30

70
100

non-origin

n side of 60.

the line The

region in the $P(60, 0)$

80 90
X

feasible region determined by the system of constraints from (ii) to (vi). The feasible region is bounded.

Step III. The coordinates of the corner point P are (60, 0). Corner point

Q: It is the point of intersection of bounding lines $x = 60$ and $x + y = 100$

Putting $x = 60$, $60 + y = 100 \Rightarrow y = 100 - 60 = 40 \therefore$ Corner point Q is (60, 40).

Corner point R: It is the point of intersection of bounding lines $y = 50$ and $x + y = 100$

Putting $y = 50$, $x + 50 = 100 \Rightarrow x = 100 - 50 = 50 \therefore$ Corner

point R is (50, 50).

Corner point S: It is the point of intersection of bounding lines $y = 50$ and $x + y = 60$

Putting $y = 50$, $x + 50 = 60 \Rightarrow x = 10$

\therefore Corner point S is (10, 50).

Step IV. Now, we evaluate Z at each corner point. . 50 Class 12

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Corner Point Z = $(5x + 3y + 820)$

P(60, 0) 560

Q(60, 40) 620

R(50, 50) 610

S(10, 50) 510 = m \leftarrow Minimum By Corner Point Method,

minimum Z = 510 at (10, 50). Hence, the transportation cost is minimum, equal to ` 510, when the supplies are transported as under:

From / To D E F

A 10 50 40 \therefore 50 0 0
 $x = 10, y = 50$) B

7. An oil company has two depots A and B with capacities of 7000 L and 4000 L respectively. The company is to supply oil to three petrol pumps, D, E and F whose requirements are 4500L, 3000L and 3500L respectively. The distances (in km) between the depots and the petrol pumps is given in the following table:

Distance (in km.)

From / To A B D 3

E 6 4

F 3 2

Assuming that the transportation cost of 10 litres of oil is ` 1 per km, how should the delivery be scheduled in order that the transportation cost is minimum? What is the minimum cost?

Sol. Step I. Mathematical formulation of L.P.P.

Let x L and y L of oil be transported from depot A to petrol pumps D and E respectively. Then $\{7000 - (x + y)\}$ L will be transported to petrol pump F.

Clearly, $x \geq 0$, $y \geq 0$ and $7000 - x - y \geq 0$ ($\Rightarrow 7000 \geq x + y$) \therefore Amounts

of petrols (in litres) can't be negative)

i.e., $x \geq 0$, $y \geq 0$ and x

7 km 6 km 3 km

$+ y \leq 7000$

Now, the

requirement of

petrol pump D is

Depot

A $_{7000}$ L

x L

$(7000 - x)$ L

4500 L. Since, x L

are transported

D

y L

E

y

F

P.P. P.P.

from depot A, the P.P. 4500 L 3000 L
remaining $(4500 - (3000 - y))$ L 4 km

3500 L + - 3500) L

x) L need to be transported x
from depot B. Similarly,
 $(3000 - y)$ L and
 3 km 2 km $(4500 -)$ L x

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$3500 - (7000 - x - y) = (x +$

y -

3500) B

4000 L Depot

(

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need to be transported from depot B to petrol pumps E and F
respectively.

Clearly, $4500 - x \geq 0$, $3000 - y \geq 0$ (i.e., $4500 \geq x$, $3000 \geq y$) and $x + y$
 $- 3500 \geq 0$

i.e., $x \leq 4500$, $y \leq 3000$, $x + y \geq 3500$

Cost of transportation of 10 litres of oil is ` 1 per km \Rightarrow Cost of

transportation of 1 litre of oil is ` per km. Total transportation
cost Z is given by

$$Z = [7x + 6y + 3(7000 - x - y) + 3(4500 - x) + 4(3000 - y) + 2(x + y - 3500)]$$

$$= (3x + y + 39500)$$

We have the following mathematical model for the given problem:

Minimise $Z = (3x + y + 39500)$... (i) subject to the constraints: x
 ≥ 0 , $y \geq 0$... (ii), $x + y \leq 7000$... (iii), $x \leq 4500$... (iv) $y \leq 3000$... (v),
 $x + y \geq 3500$... (vi)

Step II. Step II of this question Q. No. 7 is very similar to step II of
Q. No. 6 and is being left as an exercise for the reader. The reader
after drawing his or her graphs and regions should

compare shaded

with the 5000 $x = 4500$

adjoining figure. y

The

7000 (0, 7000)

feasible
region
determined by

$$\begin{aligned} x &+ y = 7000 \\ x &\geq 0, y \geq 0 \end{aligned}$$

region in the
figure is the

region is

the system of
constraints
(0, 3500)

3000 2000
S(4000, 3000)

1000
T(500, 3000) $y = 3000$
R(4500, 2500) $x + y = 7000$

from (ii) to (vi).

The feasible

bounded.

Q(4500, 0)

corner

Step III. The

(7000, 0)

points P and

P(3500, 0) 5000

X

Q are (3500, 0)

6000

X O

coordinates of the

1000 2000

and (4500, 0) respectively.

Corner point R: It is the point of intersection of bounding lines $x = 4500$ and $x + y = 7000$

Putting $x = 4500$, $4500 + y = 7000 \Rightarrow y = 7000 - 4500 = 2500$

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\therefore Corner point R is (4500, 2500).

Similarly corner points S and T are (4000, 3000) and (500, 3000) respectively.

(This is being left as an exercise for the reader).

Step IV. Now, we evaluate Z at each corner point.

Corner Point Z = $(3x + y + 39500)$

P(3500, 0) 5000

Q(4500, 0) 5300

R(4500, 2500) 5550

S(4000, 3000) 5450

T(500, 3000) 4400 = m \leftarrow Minimum By Corner Point Method,

minimum Z = 4400 at (500, 3000). Hence, the transportation cost is minimum, equal to 4400, when the supplies are transported as under:
From / To D E F

A 500L 3000L $x = 500$,

3500L \therefore

B 4000L 0L 0L $y = 3000$

8. A fruit grower can use two types of fertilizer in his garden, brand P and brand Q. The amounts (in kg) of nitrogen, phosphoric acid, potash, and chlorine in a bag of each brand are given in the table. Tests indicate that the garden needs at least 240 kg of phosphoric acid, at least 270 kg of potash and at most 310 kg of chlorine.

	kg per bag	
	Brand P	Brand Q
Nitrogen	3	3.5
Phosphoric acid	1	2
Potash	3	1.5
Chlorine	1.5	2

If the grower wants to minimise the amount of nitrogen added to the garden, how many bags of each brand should be used? What is the minimum amount of nitrogen added in the garden?

Sol. Step I. Mathematical formulation of L.P.P.

Let the fruit grower use x bags of brand P and y bags of brand Q. The given data is condensed in the following table.

	fertilizer of bags		Nitrogen	Phosphoric	Potash	Chlorine	Acid
P	x	3	1	3	1.5	Q	y
Q	3.5	2	1.5	2	Amount of nitrogen	$= 3x + 3.5y$	
Let Z	$= 3x + 3.5y$						

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We have the following mathematical model for the given problem:

Minimise $Z = 3x + 3.5y$... (i) subject to the constraints:

$$x + 2y \geq 240 \text{ (Phosphoric acid constraint) ... (ii)}$$

[Given: The garden needs at least (i.e., \geq) 240 kg of phosphoric acid]

$$3x + 1.5y \geq 270 \text{ or } 3x + y \geq 270$$

[Given: The garden atleast 270 kg of potash]

Dividing by 3 and multiplying by 2,

$$\text{or } 2x + y \geq 180 \text{ (Potash constraint) ... (iii)}$$

$$1.5x + 2y \leq 310 \text{ or } x + 2y \leq 310$$

[Given: The garden needs at the most i.e., \leq 310 kg of chlorine]

Multiplying by 2, $3x + 4y \leq 620$.

$$\text{or } 3x + 4y \leq 620 \text{ (Chlorine constraint) ... (iv) } x, y \geq 0 \text{ ... (v) } . .$$

Amounts of phosphoric acid, potash and chlorine can't be negative)

Step II. The region for constraint (v), $x, y \geq 0$

\Rightarrow Feasible region is in first quadrant.

Table of values for the line $x + 2y = 240$ of constraint (ii)

x 0 240
120 0

Let us draw the line joining the points (0, 120) and (240, 0). Let us test for origin ($x = 0, y = 0$) in constraint (ii), $x + 2y \geq 240$, we have $0 \geq 240$ which is not true.

\therefore Region for constraint (ii) is on the non-origin side of the line $x + 2y = 240$ i.e., region is half plane on the above side of the line $x + 2y = 240$. Table of values for the line $2x + y = 180$ for constraint (iii) $x = 90, y = 180$

Let us draw the line joining the points (0, 180) and (90, 0). Let us test for origin ($x = 0, y = 0$) in constraint (iii) $2x + y \geq 180$, we have $0 \geq 180$ which is not true.

\therefore Again region for constraint (iii) is also the half-plane not containing the origin i.e., on the non-origin side of the line $2x + y = 180$. Table of

values for the line $3x + 4y = 620$ for constraint (iv)

$$x = 200.7$$

$y = 155$

Let us draw the line joining the points (0, 155) and (200.7, 0). Let us test for origin ($x = 0, y = 0$) in $3x + 4y \leq 620$, we have $0 \leq 620$ which is true.

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\therefore Region for constraint (iv) is on the origin side of the line $3x + 4y = 620$.

The shaded region in the figure is the feasible region determined by the system of constraints from (ii) to (v). The feasible region is bounded.

Y

(0, 180)

(0, 155)

140

(0, 120) 100

80

60

C(20, 140) B(40, 100)

2

x

A(140, 50)

40

+

y

3

x

$x + 2 = 24$

$0 = 180$

y

20

$0 = 180$

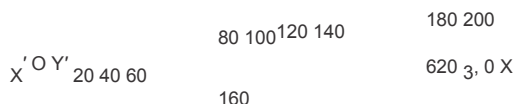
0

(90, 0)

$+ 4 = 620$

0 y

(240, 0)



Step III. Let us find the corner points A, B and C. Corner point A:

It is the point of intersection of the lines $x + 2y = 240$ and $3x + 4y = 620$

$$4y = 620$$

Second Eqn. $- 3 \times$ First equation gives

$$3x + 4y - 3x - 6y = 620 - 720$$

$$\Rightarrow -2y = -100 \Rightarrow y =$$

—

$$= 50$$

Putting $y = 50$ in $x + 2y = 240$, we

have

$$x + 100 = 240 \Rightarrow x =$$

$$140$$

\therefore Corner point A

is $(140, 50)$.

Corner point B: It is the point of intersection of bounding lines $x + 2y = 240$ and $2x + y = 180$

First Eqn. $- 2 \times$ Second equation gives

$$x + 2y - 4x - 2y = 240 - 360$$

$$\Rightarrow -3x = -120 \Rightarrow x = 40$$

Putting $x = 40$ in $x + 2y = 240$, we have

$$40 + 2y = 240 \Rightarrow 2y = 200 \Rightarrow y = 100$$

\therefore Corner point B is $(40, 100)$.

Corner point C: It is the point of intersection of bounding lines $2x + y = 180$ and $3x + 4y = 620$

Second Eqn. $-4 \times$ First equation gives

$3x - 8x = 620 - 720 \Rightarrow -5x = -100 \Rightarrow x = 20$ Putting $x = 20$ in $2x + y = 180$, we have $40 + y = 180 \Rightarrow y = 140 \therefore$ Corner point C is (20, 140).

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Step IV. Now, we evaluate Z at each corner point.

Corner Point $Z = 3x + 3.5y$

A(140, 50) 595

B(40, 100) 470 = m \leftarrow Minimum

C(20, 140) 550

By Corner Point Method, minimum $Z = 470$ at (40, 100). \therefore Minimum amount of nitrogen = 470 kg when 40 bags of brand P and 100 bags of brand Q are used.

9. Refer to Question 8. If the grower wants to maximise the amount of nitrogen added to the garden, how many bags of each brand should be added? What is the maximum amount of nitrogen added?

Sol. From the above Table of Step IV in solution of question 8, we find that $Z = 595$ is maximum at (140, 50).

\therefore Maximum amount of nitrogen = 595 kg when 140 bags of brand P and 50 bags of brand Q are used.

10. A toy company manufactures two types of dolls, A and B. Market tests and available resources have indicated that the combined production level should not exceed 1200 dolls per week and the demand for dolls of type B is at most half of that for dolls of type A. Further, the production level of dolls of type A can exceed three times the production of dolls of other type by at most 600 units. If the company makes profit of ` 12 and ` 16 per doll respectively on dolls A and B, how many of each should be produced weekly in order to maximise the profit? Solve it graphically.

Sol. Step I. Mathematical Formulation of L.P.P.

Let x dolls of type A and y dolls of type B be produced to have the maximum profit.

Given: Company makes profit of ` 12 and ` 16 per doll respectively on doll A and B.

\Rightarrow Objective function is

Profit $Z = 12x + 16y$

Constraint on number of dolls

Given: Combined production level of dolls should not exceed 1200 dolls per day.

$\Rightarrow x + y \leq 1200$... (i) Again given demand for dolls of type B is at

most half that for dolls of type A. At most $\Rightarrow y \leq \frac{1}{2}x$... (ii) Again given: production level of dolls of type A can exceed three times the production of dolls of other type (B) by at most 600 units.

$\Rightarrow x \leq 3y + 600$

$\Rightarrow x - 3y \leq 600$... (iii) Also $x \geq 0, y \geq 0$ because number of dolls can't be negative. **Step II.** To draw the graphs for regions of all constraints and

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locate the common feasible region.

Constraint (i) is $x + y \leq 1200$

Replacing \leq by $=$, $x + y = 1200$

$$x \ 0 \ 1200 \ (0, 1200)$$

$$y \ 1200 \ 0 \ (1200, 0)$$

\therefore Graph of $x + y = 1200$ is the straight line joining the points $(0, 1200)$ and $(1200, 0)$.

Let us test for origin in (i),

Put $x = 0$ and $y = 0$ in (i), $0 \leq 1200$ which is true. \therefore Region given by (i) is towards the origin and is being shown by horizontal lines. Constraint

(ii) is $y \leq$

Let us draw graph of $y =$

$$x \ 0 \ 400$$

$$y \ 0 \ 200$$

\therefore Graph of $y =$ is the straight line joining $(0, 0)$ and $(400, 200)$. Let us test for $(1200, 0)$ in (ii), $0 \leq 600$ which is true. \therefore Region given by (ii) is towards the point $(1200, 0)$, shown by vertical lines.

Constraint (iii) is $x - 3y \leq 600$

Let us draw the graph of $x - 3y = 600$

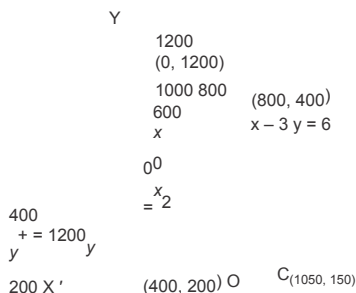
$$x \ 0 \ 600$$

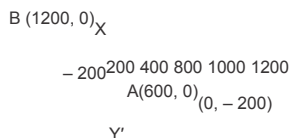
$$y \ -200 \ 0$$

\therefore Graph of $x - 3y = 600$ is the straight line joining the points $(0, -200)$ and $(600, 0)$.

Let us test for origin $(0, 0)$ in (iii).

Put $x = 0$ and $y = 0$ in (iii). $0 \leq 600$ which is true. \therefore Region given by (iii) is towards the origin shown by slanting lines.





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The common feasible region is bounded by quadrilateral OABC. Step III. The vertices of this feasible region are

O(0, 0); A(600, 0)

B, point of intersection of the lines:

$$x - 3y = 600$$

and $x + y = 1200$

Subtracting

$$-4y = -600$$

$$\therefore y =$$

$$= 150$$

Putting $y = 150$ in $x + y = 1200$,

$$x + 150 = 1200$$

$$\Rightarrow x = 1200 - 150 = 1050$$

\therefore Corner point B(1050, 150)

Corner point C is point of intersection of

lines: $y =$

and $x + y = 1200$

Solving $x + y = 1200 \Rightarrow 2x + x = 2400$

\therefore Corner point C is (800,

$$\Rightarrow 3x = 2400 \Rightarrow \therefore y = 400)$$

$$= 800$$

$$= 400$$

Step IV. Values of objective (profit) function Z at corner points are:

Corner point	Value of objective function $Z = 12x + 16y$
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O(0, 0)	$Z = 12(0) + 16(0) = 0$
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A(600, 0)	$Z = 12(600) + 16(0) = 7200$
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	$Z = 12(1050) + 16(150)$
--	--------------------------

B(1050, 150)	
$= 12600 + 2400 = 15000$	$Z = 16000 \rightarrow M$
	$= 12(800) + 16(400)$

C(800, 400)	$= 9600 + 6400$
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\therefore Maximum profit is ` 16000 when $x = 800$, $y = 400$. . 58